

Do as I Say not as I Do
Stealth Modification of
Programmable Logic Controllers
I/O by Pin Control Attack

ALI ABBASI


SYSSEC GROUP,
RUHR UNIVERSITY BOCHUM, GERMANY
& SCS GROUP
UNIVERSITY OF TWENTE, NETHERLANDS

MAJID HASHEMI

PARIS, FRANCE

Who we are

- Ali Abbasi, visiting researcher at chair of system security of Ruhr University Bochum and PhD student at Distributed and Embedded Systems Security Group, University of Twente, The Netherlands.

( @bl4ckic3)

- Majid Hashemi, R&D researcher ( @m4ji_d).

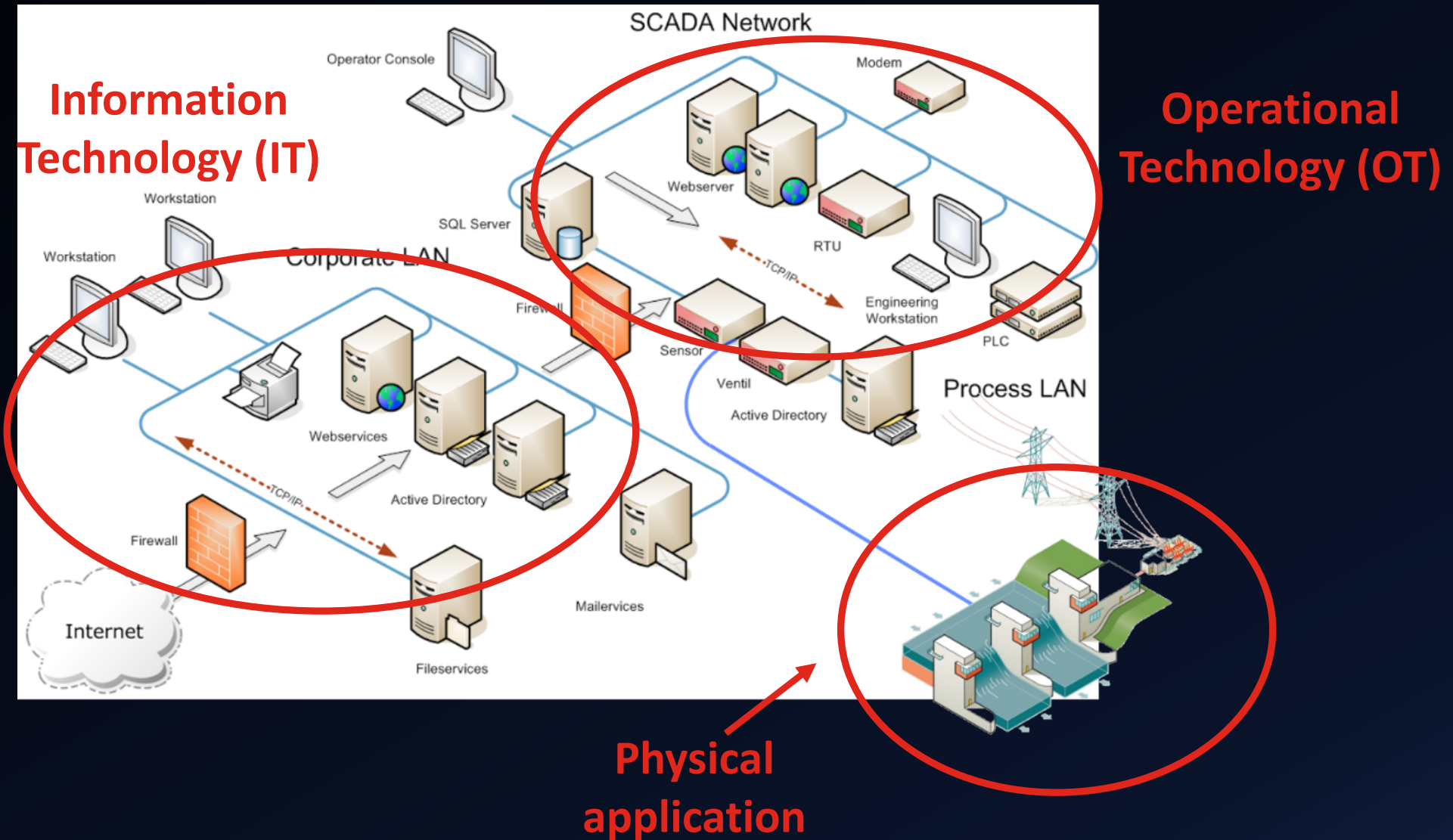
Agenda

- Background on Process Control
- Background on existing attacks and defenses for embedded systems
- Applicable Defenses for PLCs
- Background on Pin Control
- The Problem with Pin Control
- Rootkit variant
- Non-rootkit variant
- Demo
- Discussions

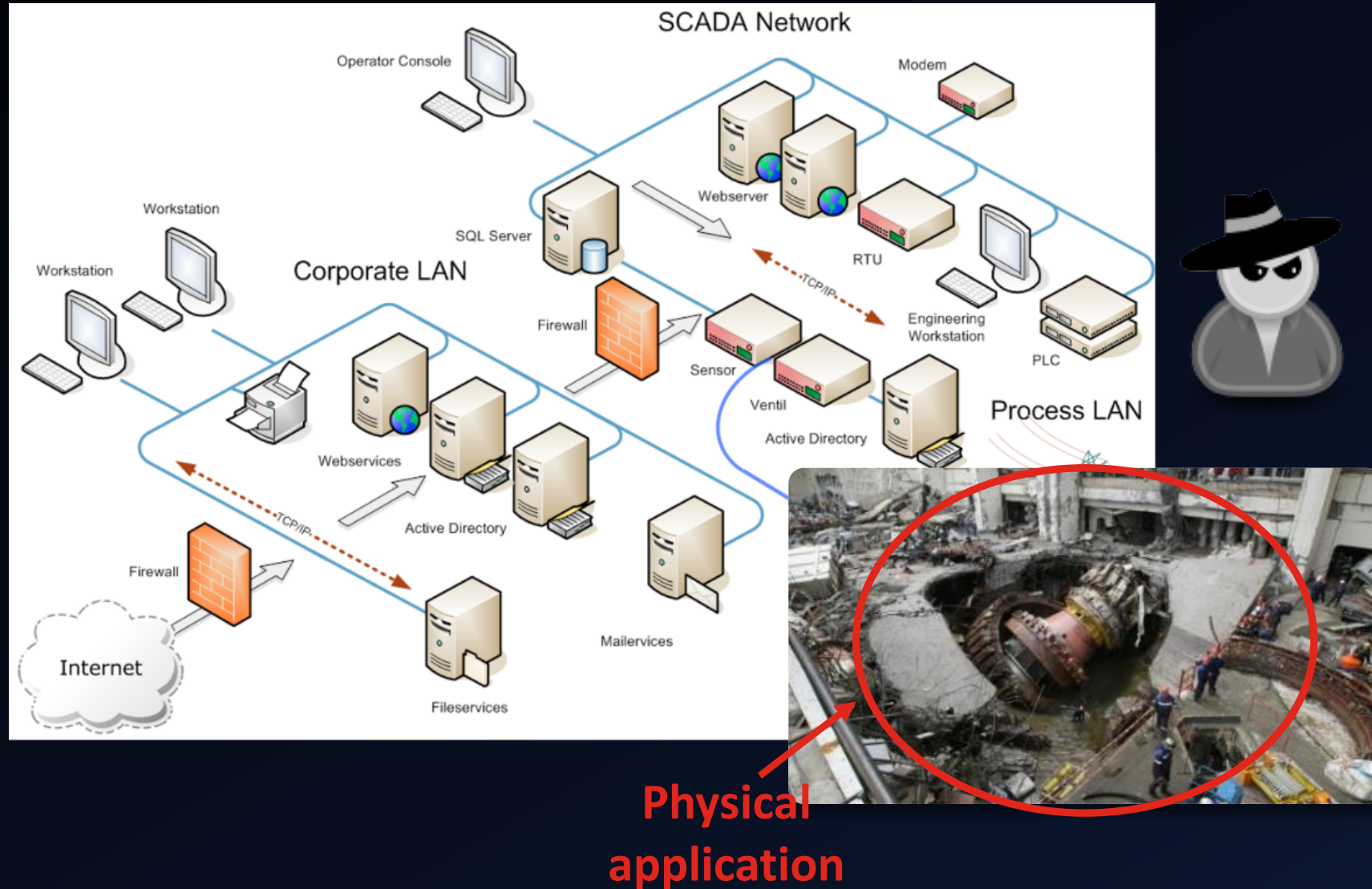
What this talk is about?

- The talk is trying to uncover existing design flaw in PLCs.
- The attack can be used in future by attackers.
- We are not unveiling fully functional malware for PLCs.
- No exploitation techniques, no 0day leak
- We are not going to mention any vendor name.

Industrial Control System



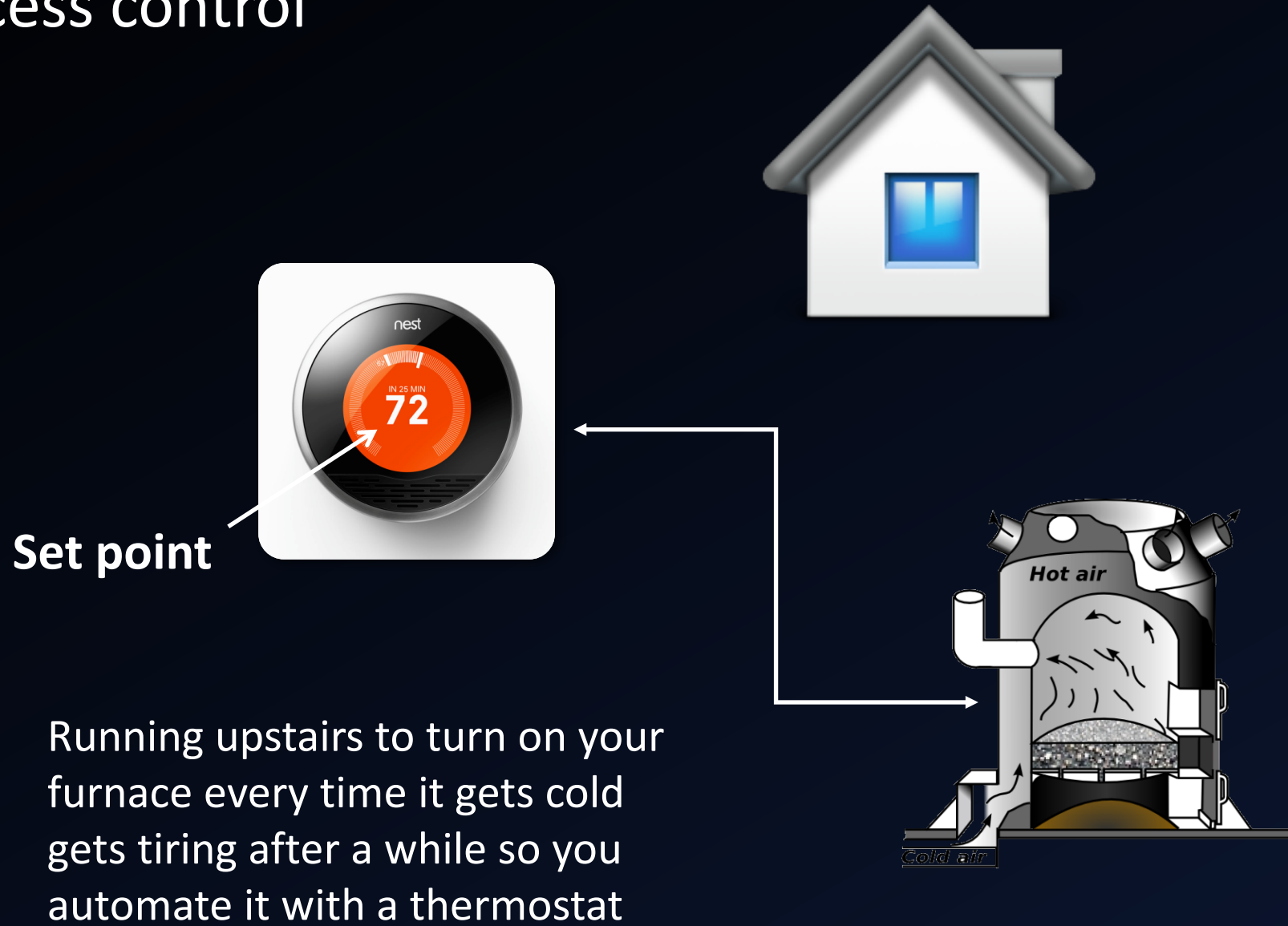
Industrial Control System hacking



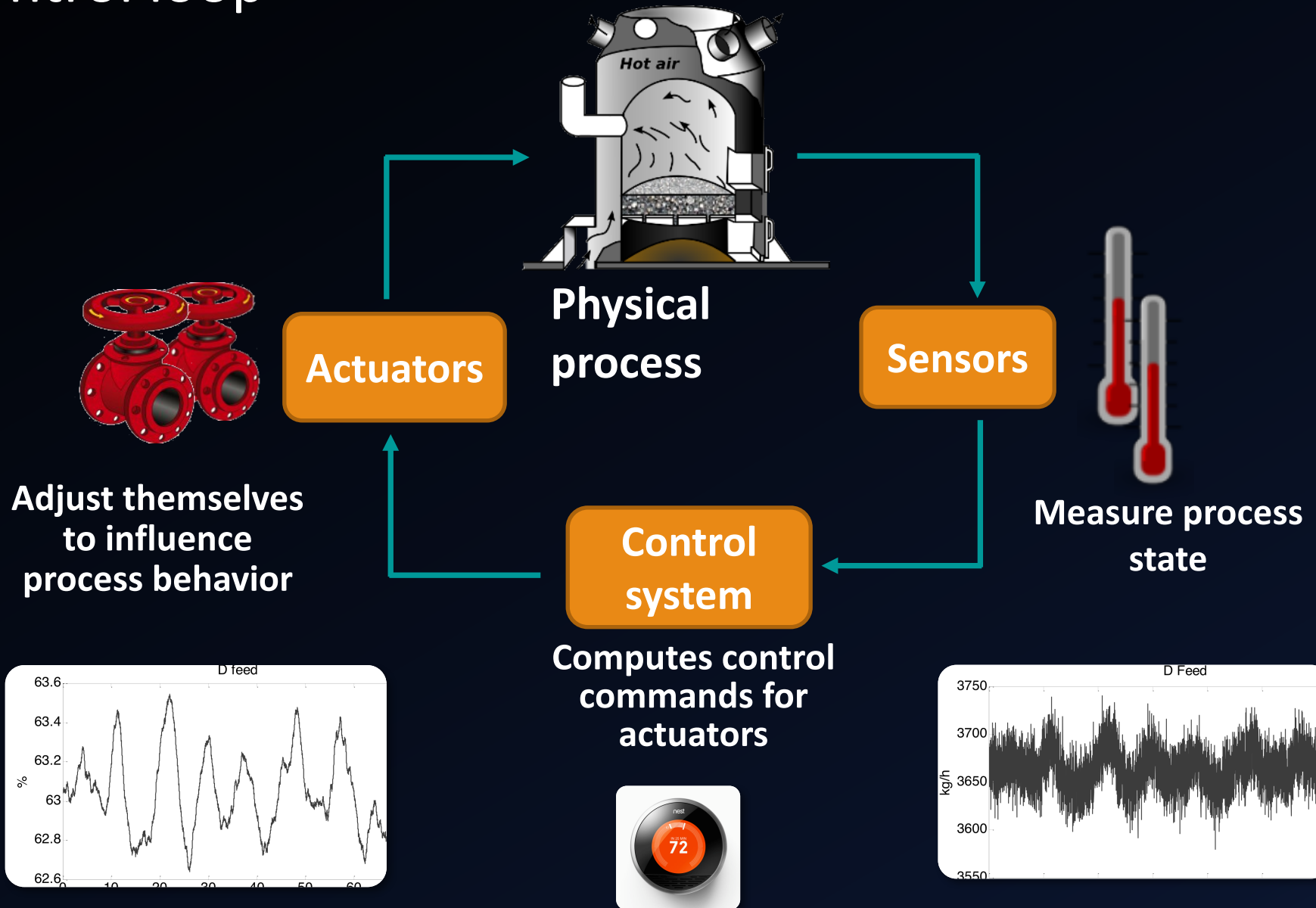


Process control 101

Process control



Control loop



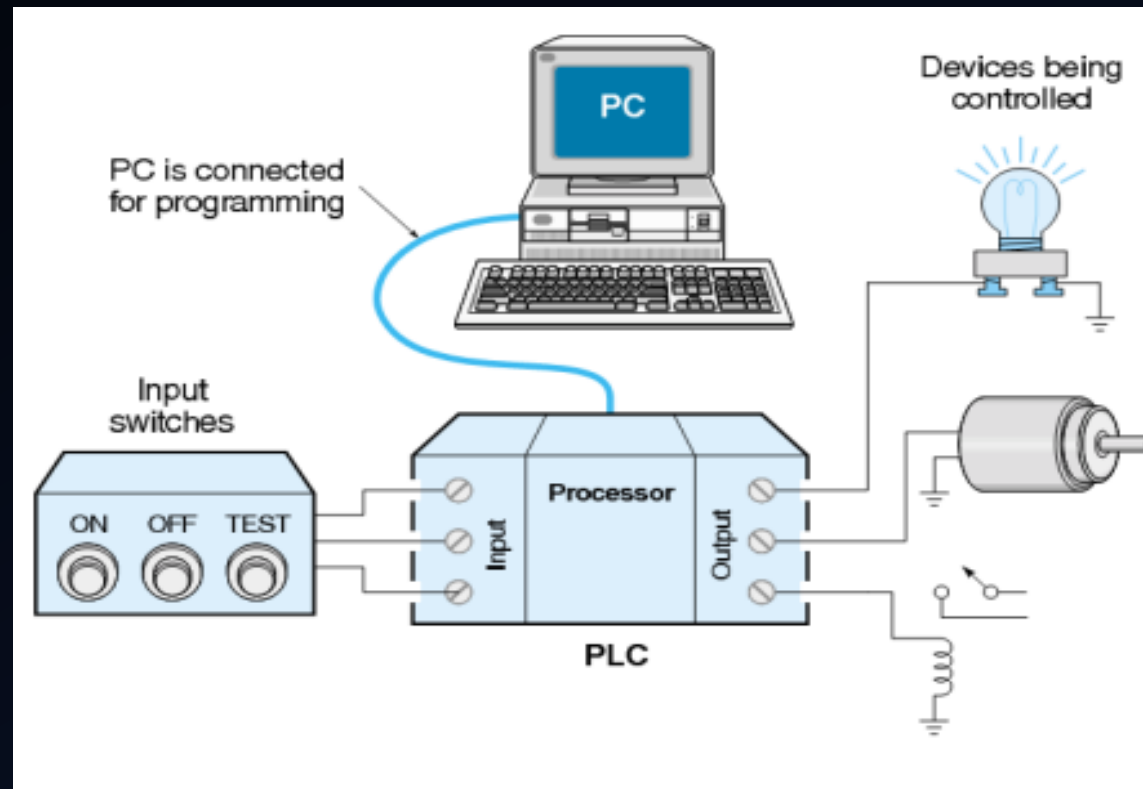
Control equipment

- In large –scale operations control logic gets more complex than a thermostat
- One would need something bigger than a thermostat to handle it
- Most of the time this is a programmable logic controller (PLC)



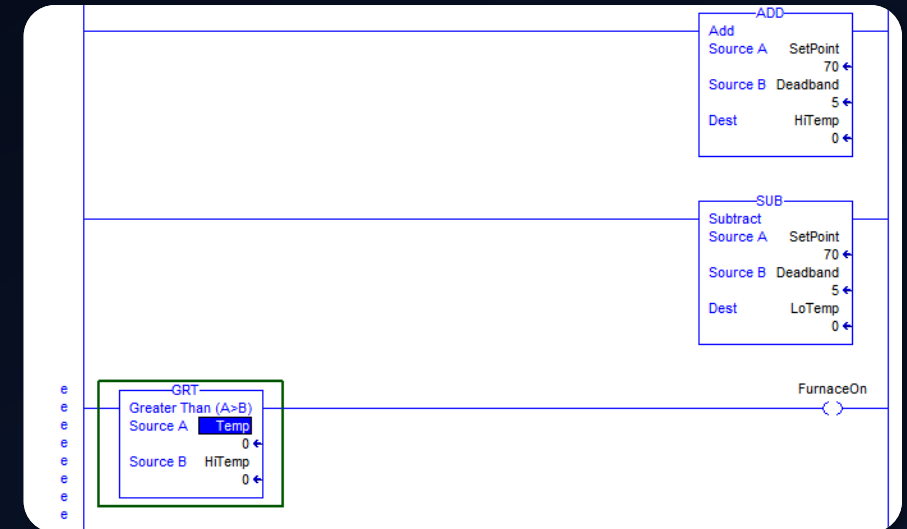
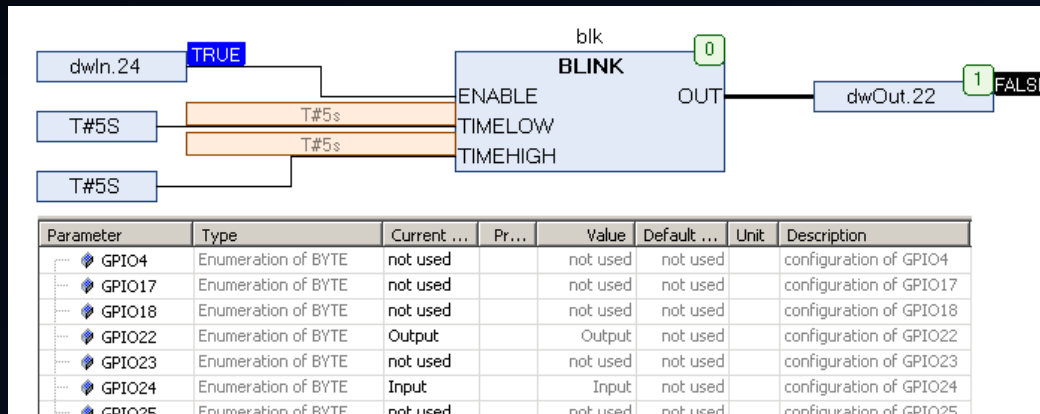
What is a PLC?

- An Embedded System with RTOS running logic.



Control logic

- It is programmed graphically most of the time
- Defines what should/should not happen
 - Under which conditions
 - At what time
 - Yes or No proposition



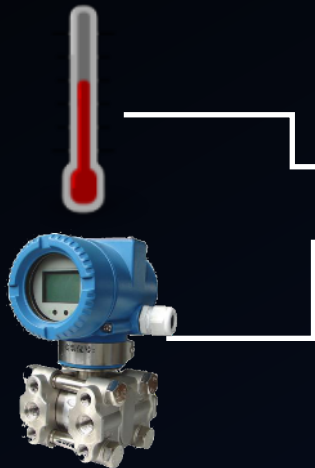
[if input 1] AND [input 2 or input 11]
-> [do something in output 6]

If tank pressure in PLC 1 > 1800
reduce inflow in PLC 3

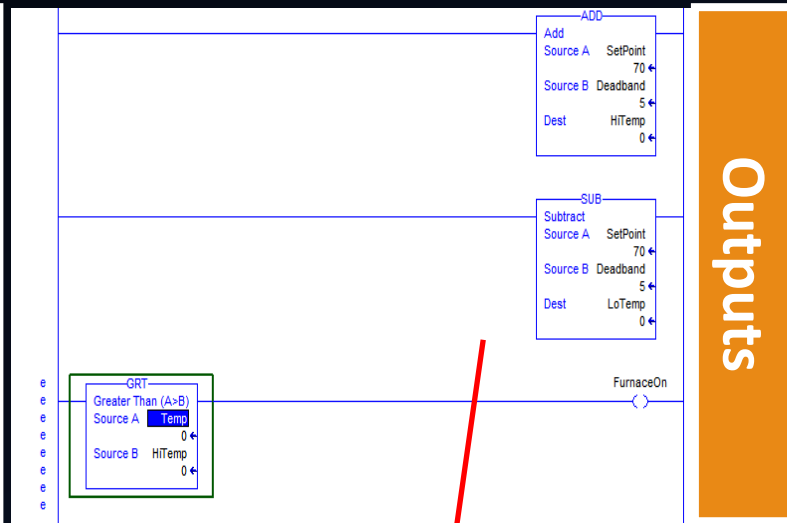
How PLC Works

1. Copy data from inputs to temporary storage
2. Run the logic
3. Copy from temporary storage to outputs

Sensors

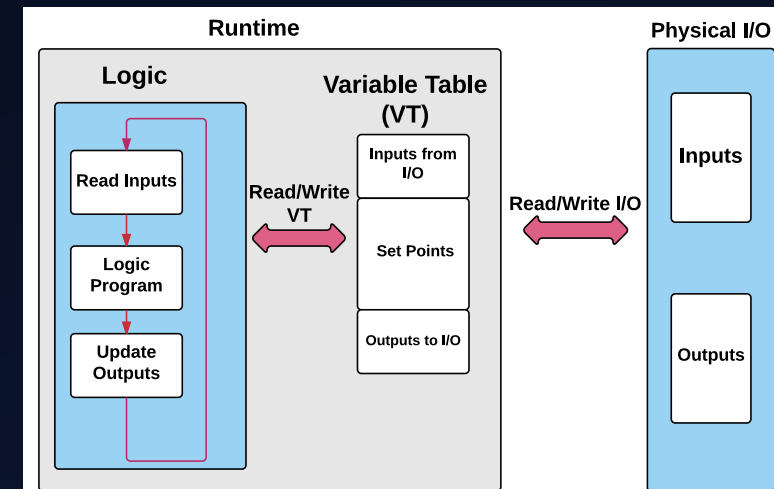


Inputs



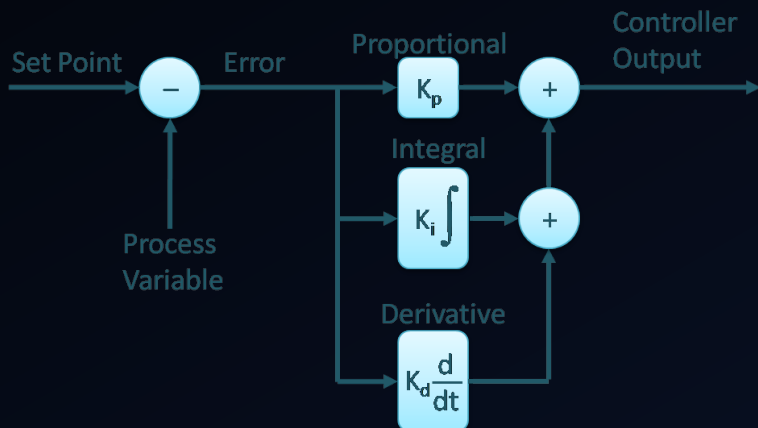
Outputs

Actuators

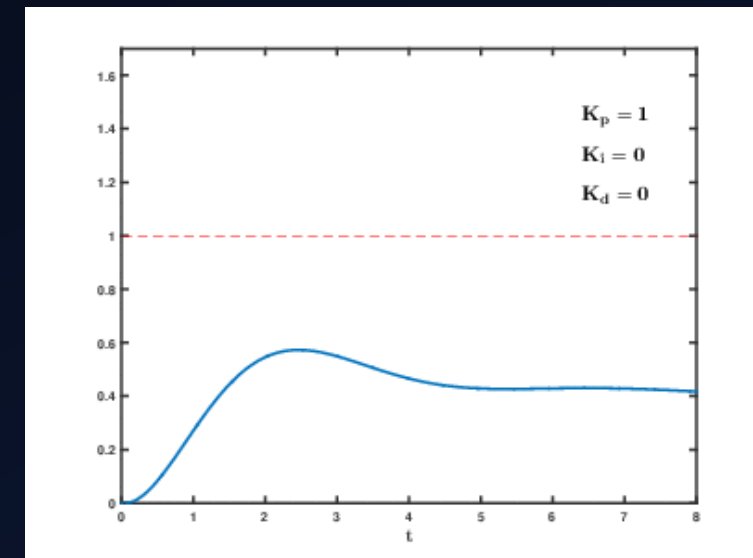
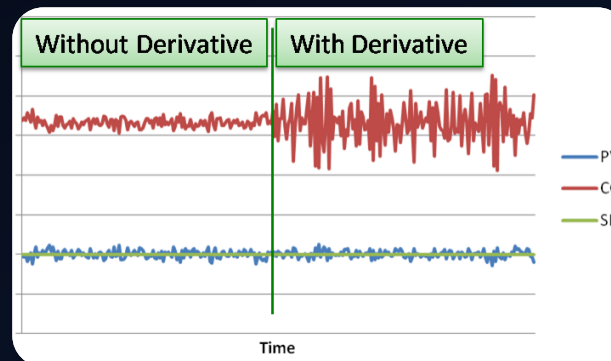


Control algorithm

- Used to compute output based on inputs received from control logic
- **PID: proportional, integral, derivative** – most widely used control algorithm on the planet
- PI controllers are most often used



$$u(t) = K \left(e(t) + \frac{1}{T_i} \int_0^t e(\tau) d\tau + T_d \frac{de(t)}{dt} \right)$$



Current attacks against embedded systems

- Authentication bypass
 - Attacker find a backdoor password in the PLC.
- Firmware modification attacks
 - Attacker upload new firmware to the PLC
- Configuration manipulation attacks
 - Attacker modify the logic
- Control Flow attacks
 - Attacker find a buffer overflow or RCE in the PLC
- Hooking functions for ICS malwares

Current defenses for embedded systems

- Attestation
 - memory attestation
- Firmware integrity verification
 - Verify the integrity of firmware before its being uploaded
- Hook detection
 - Code hooking detection
 - Detect code hooking
 - Data hooking detection
 - Detect data hooking

Requirement for Applicable Defenses for PLCs

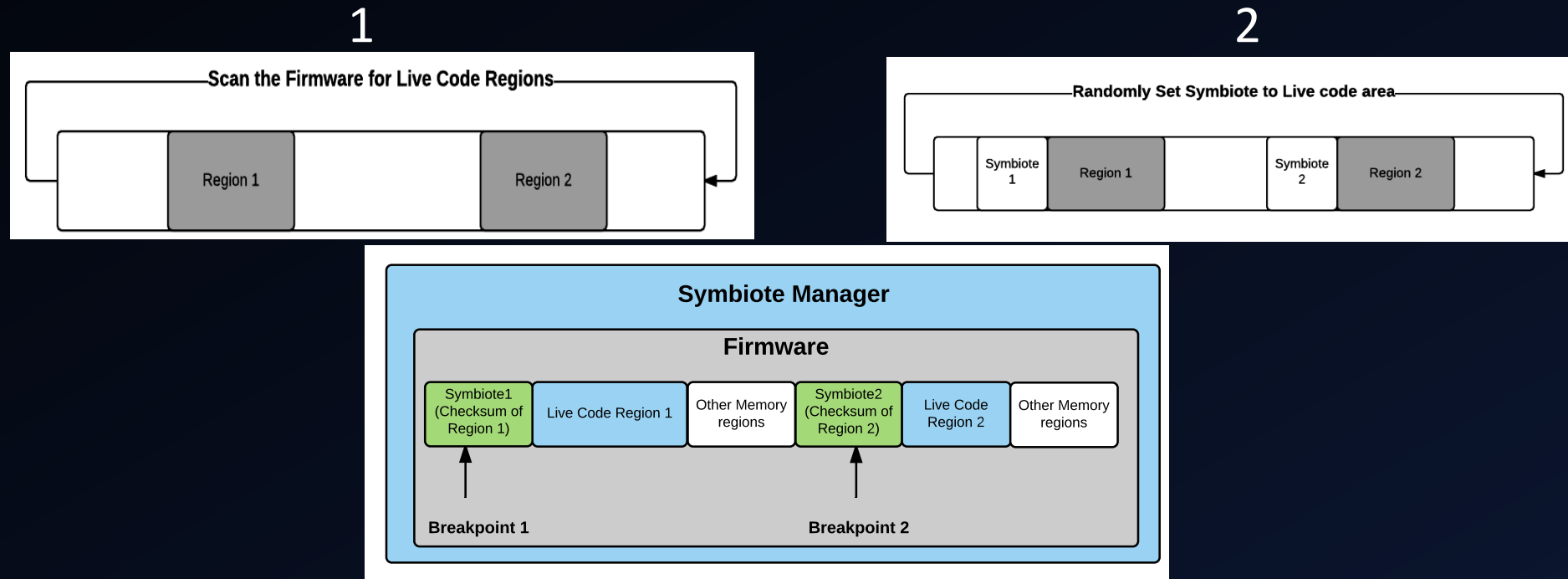
- Designed for embedded devices running modern OS.
- No hardware modifications.
- Limited CPU overhead.
- No virtualization.

System-level protection for PLCs

- Trivial Defenses:
 - Logic Checksum
 - Firmware integrity verification
- Non-trivial software-based HIDS applicable to PLCs
 - Doppelganger (Symbiote Defense): an implementation for software symbiotes for embedded devices
 - Autoscapy JR: A host based intrusion detection which is designed to detect kernel rootkits for embedded control systems

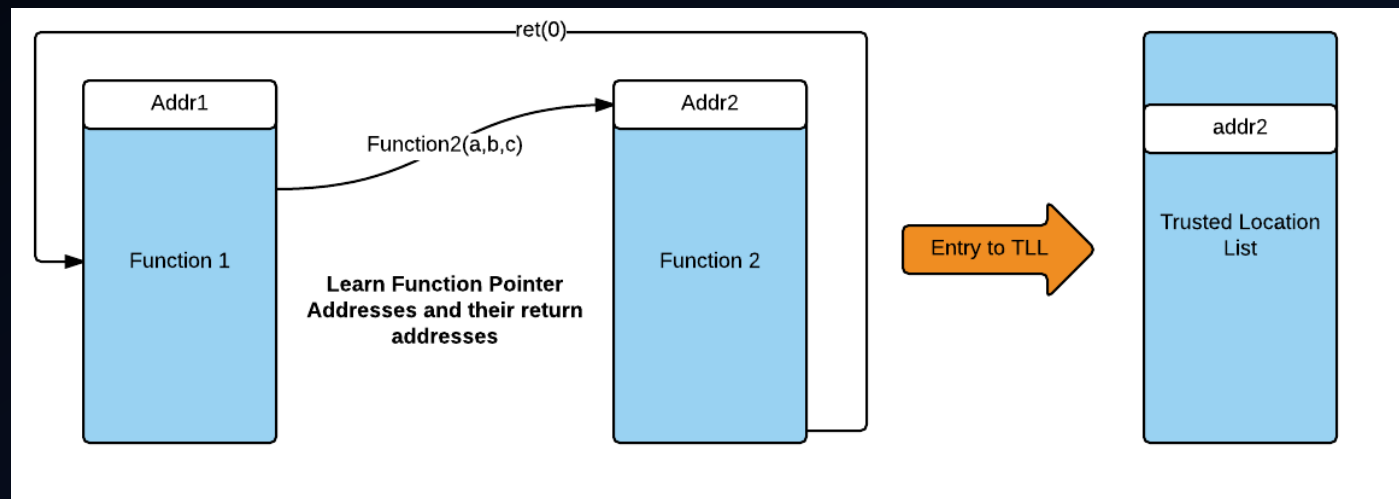
How Doppelganger Works

- Scan the firmware of the device for live code regions and insert symbiotes randomly.



How Autoscopy Jr works

- Tries to Detects function hooking by learning
- Verifies the destination function address and returns with the values and addresses in TLL (Trusted Location List)



Debug Registers

- Designed for debugging purpose.
- Function hooking intercept the function call and manipulate the function argument.
- We use debug registers in ARM processors to intercept memory access (No function interception, no function argument manipulation)

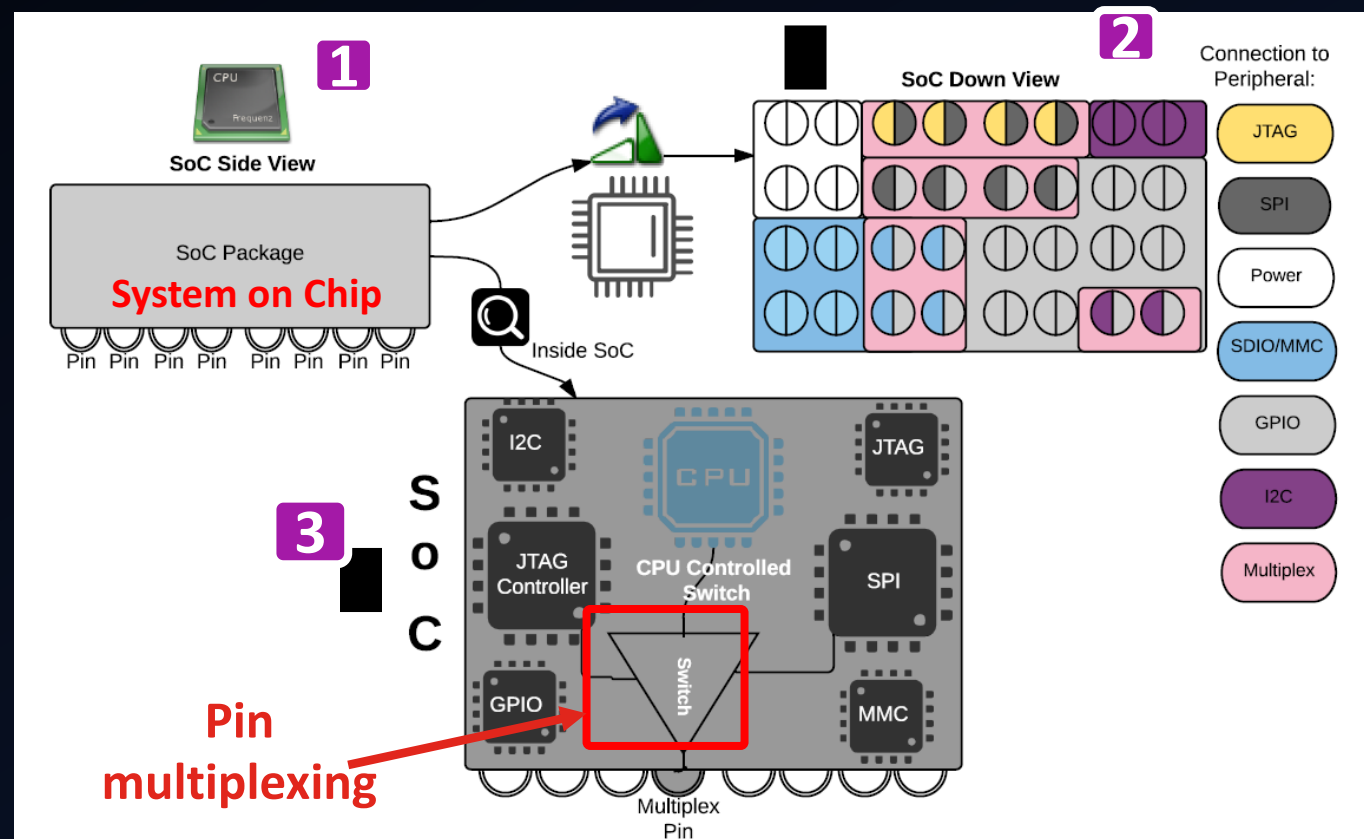


```
==Phrack Inc.==  
  
Volume 0x0c, Issue 0x41, Phile #0x08 of 0x0f  
  
|-----=[ Mistifying the debugger, ]-----|  
|-----=[   ultimate stealthness   ]-----|  
|-----=[ halfdead@phear.org ]-----|  
  
--[ Introduction  
  
Over the years, there have been a plethora of techniques and methods of  
hiding one's presence in a hacked system. Many of them were focused on  
directly tampering the system call table, others were modifying the  
interrupt handler, while others were operating at the VFS layer. But all  
of them were modifying the underlying operating system in a very visible  
manner, making them easily detected.  
  
In the article I will present a technique that is able to achieve ultimate  
stealthness in kernel rootkits, by using a common x86 feature, the  
debugging mechanism. Although it works on any IA-32 compatible platform,
```


Background on Pin Control

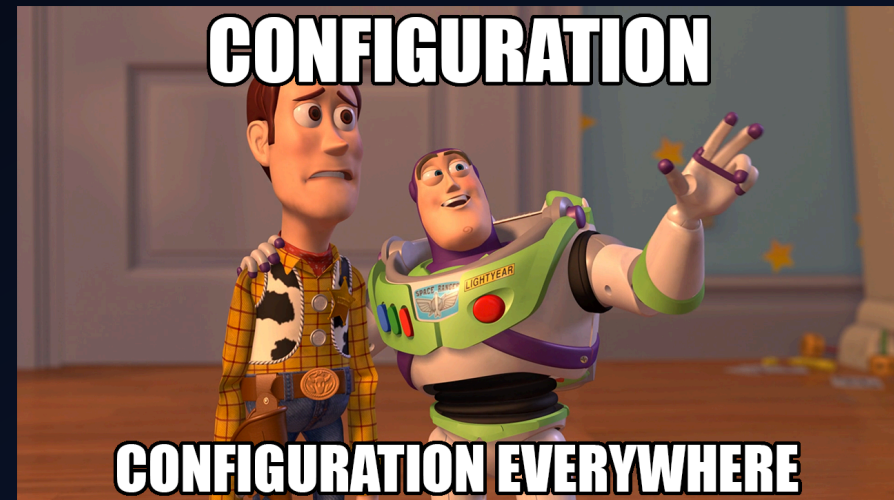
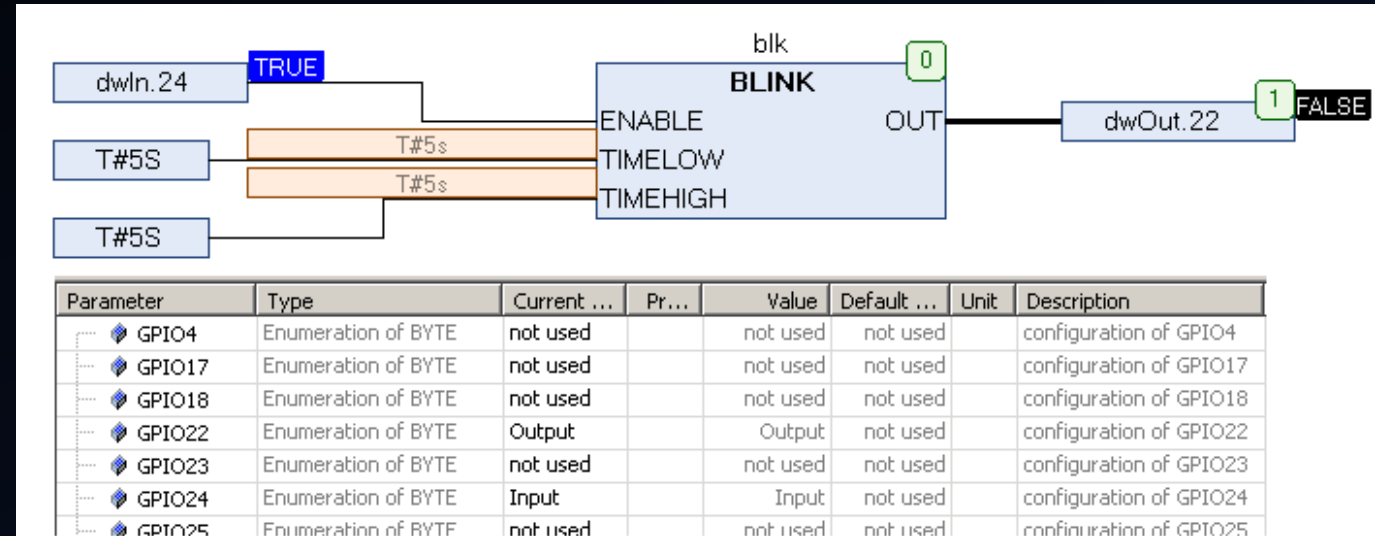
Pin Control subsystem

- Pin multiplexing (type)
- Pin configuration (in/out)

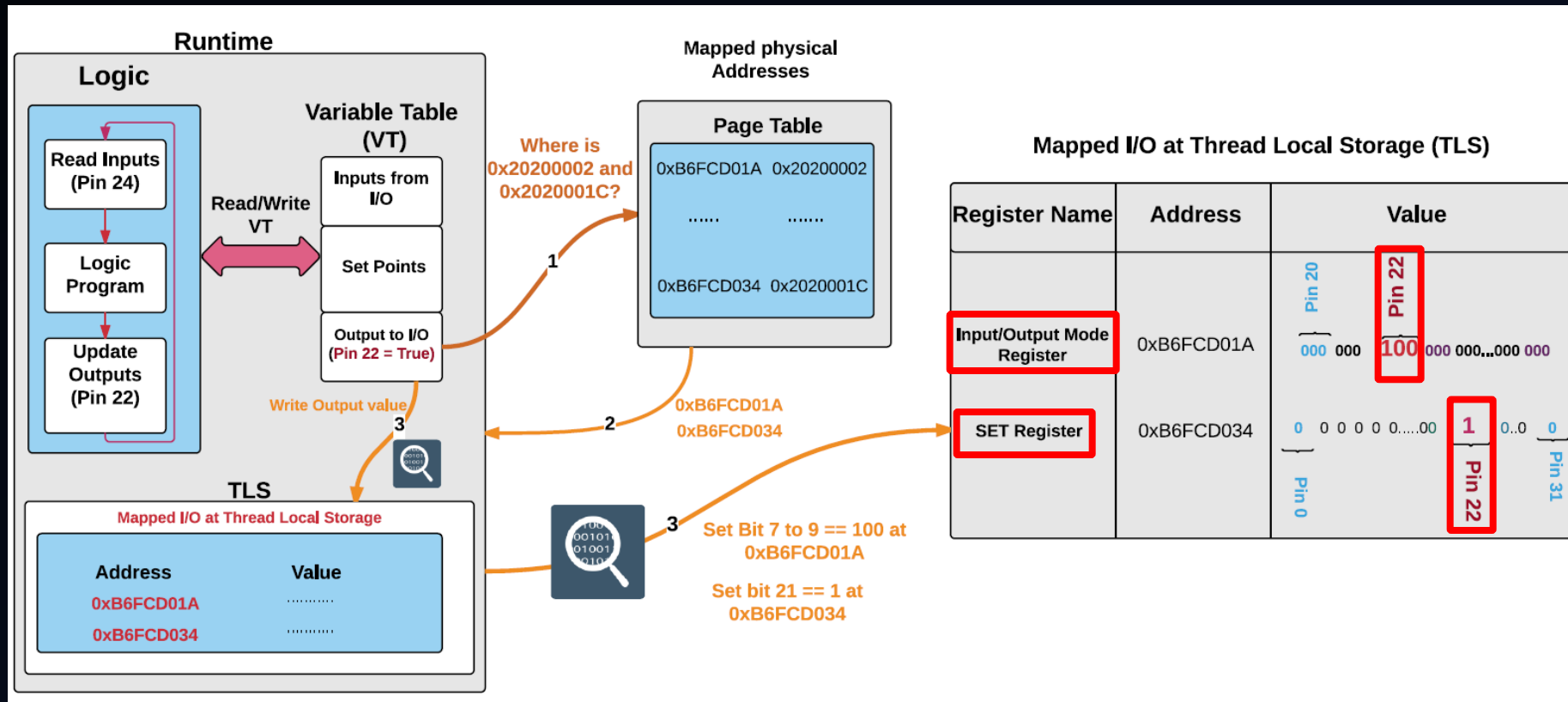


Pin Configuration

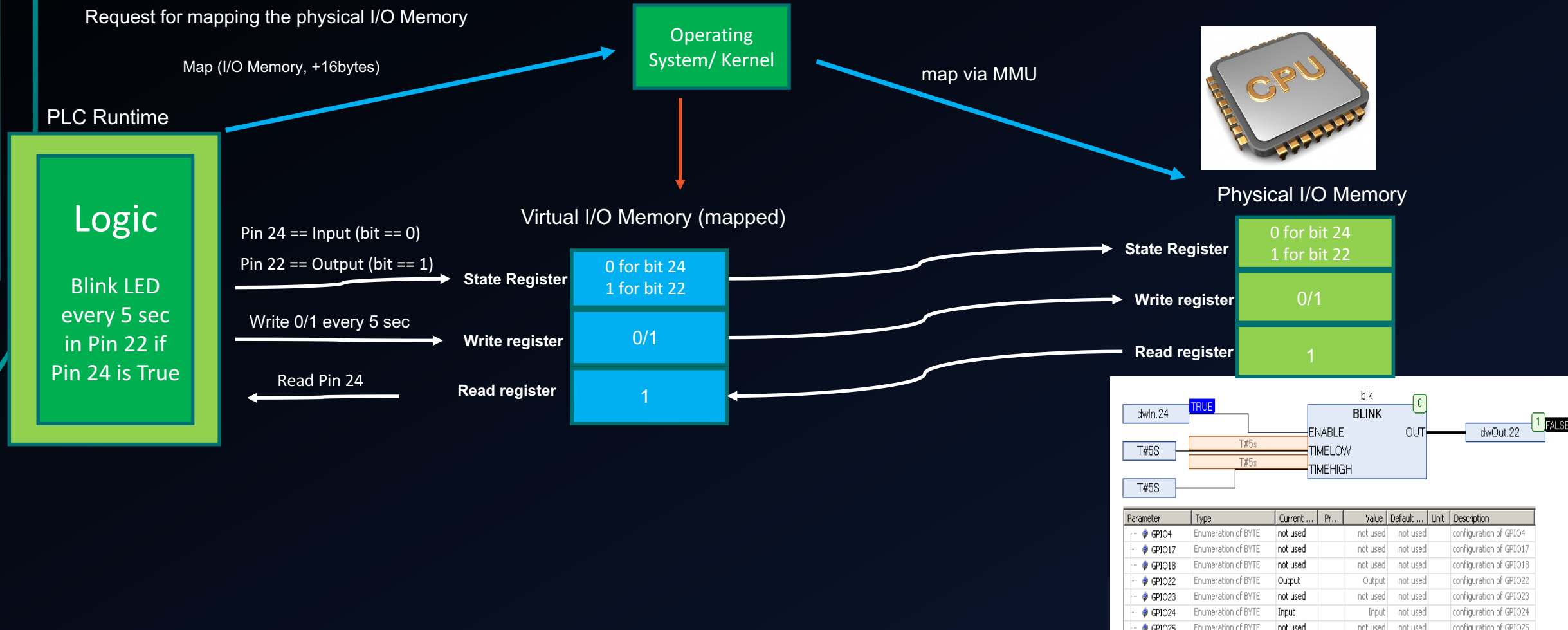
- Input Pin
 - readable but not writeable
- Output Pin
 - readable and writeable



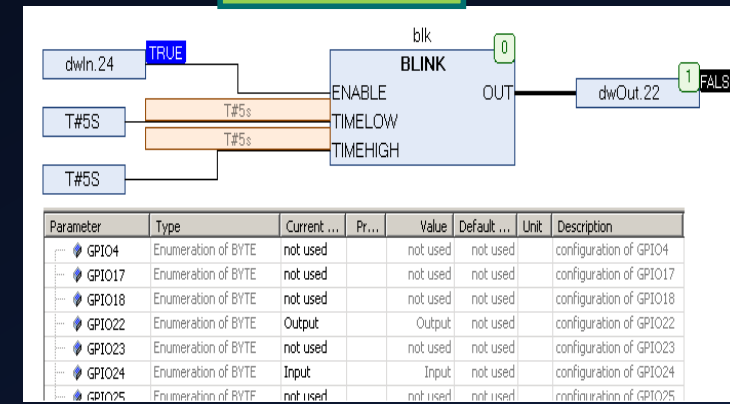
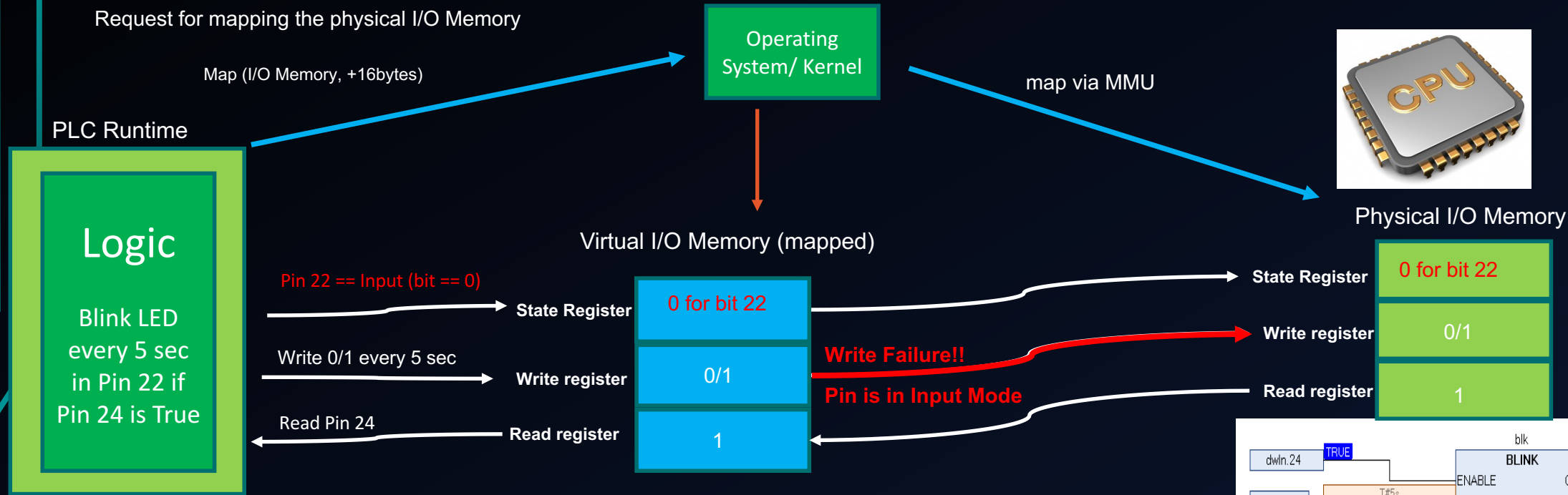
How PLC controls I/O



Introducing Pin Control Attack: A Memory Illusion



Introducing Pin Control Attack: A Memory Illusion



Think of copying files to USB drive

- Similar mapping between physical and virtual addresses
- If USB drive is removed during copy operation, OS reports a warning back



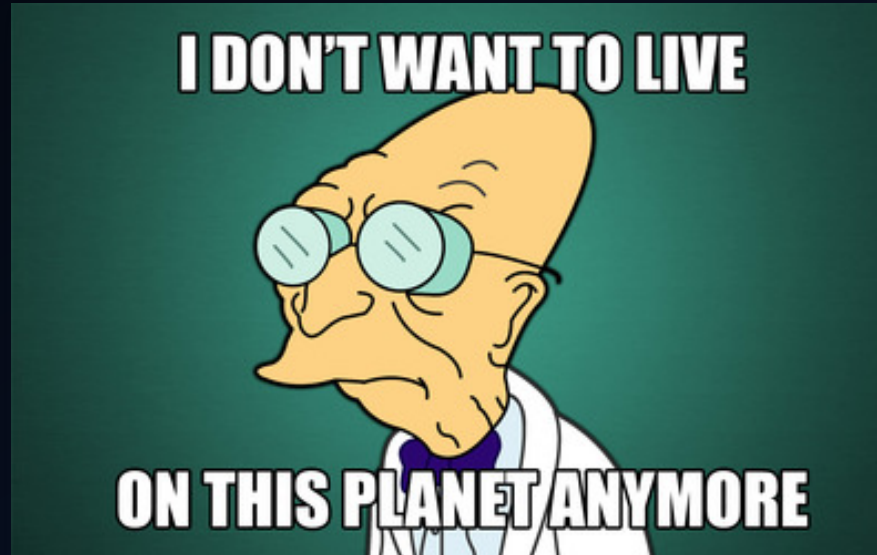
Lets look at it.

Demo 1

Digital

Nobody thought about the same issue for PLCs

- Shouldn't the PLC runtime fail or get terminated because of I/O failure?
 - Nope!



- PLC design was always about paramount reliability of real-time execution, HIGH up-time and long-term useful life in harsh environmental conditions
- Malicious manipulation of PLC were not part of design considerations :-)

Security concerns regarding pin control

- No interrupt for pin configuration
 - How the OS knows about the modification of pin configuration?
 - What if somebody modifies configuration of a pin at runtime?
 - By switching input pin into output pin, it is possible to write arbitrary value into its physical address
- No Interrupt for pin multiplexing
 - How OS knows about modification of pin multiplexing?
 - What if somebody multiplex a pin at runtime?
 - By multiplexing pin it is possible to prevent runtime from writing value into output pin



Problem statement

- What if we create an attack using pin control that:
 - Do not do function hooking
 - Do not modify executable contents of the PLC runtime.
 - Do not change the logic file
- Obviously we consider other defenses available (e.g. logic checksum is also there)

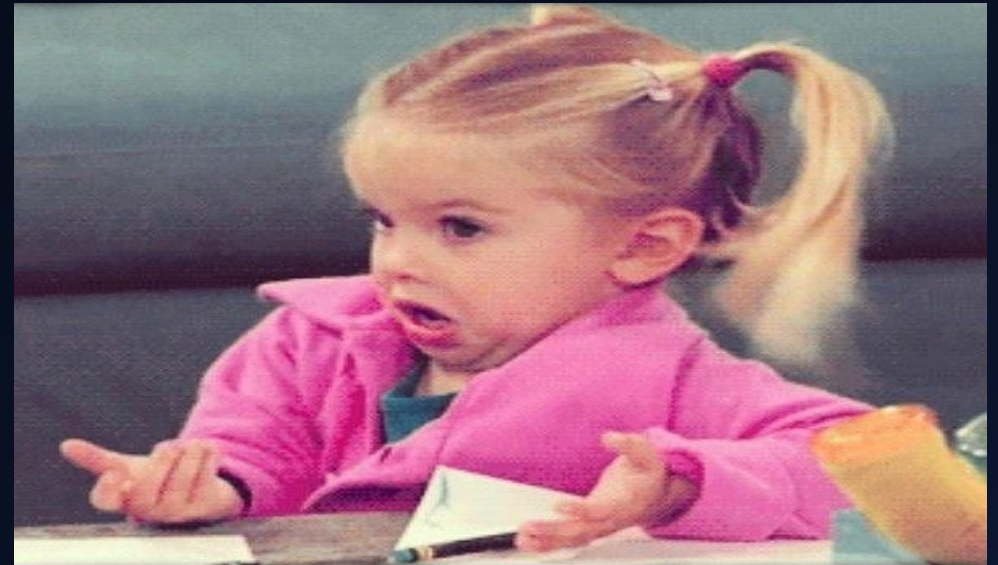




Pin Control Attack

Pin Control Attack

- Pin Control Attack:
 - manipulate the I/O configuration (Pin Configuration Attack)
 - manipulate the I/O multiplexing (Pin Multiplexing Attack)
- PLC OS will never know about it.



Two options to achieve the same



1 First version: rootkit

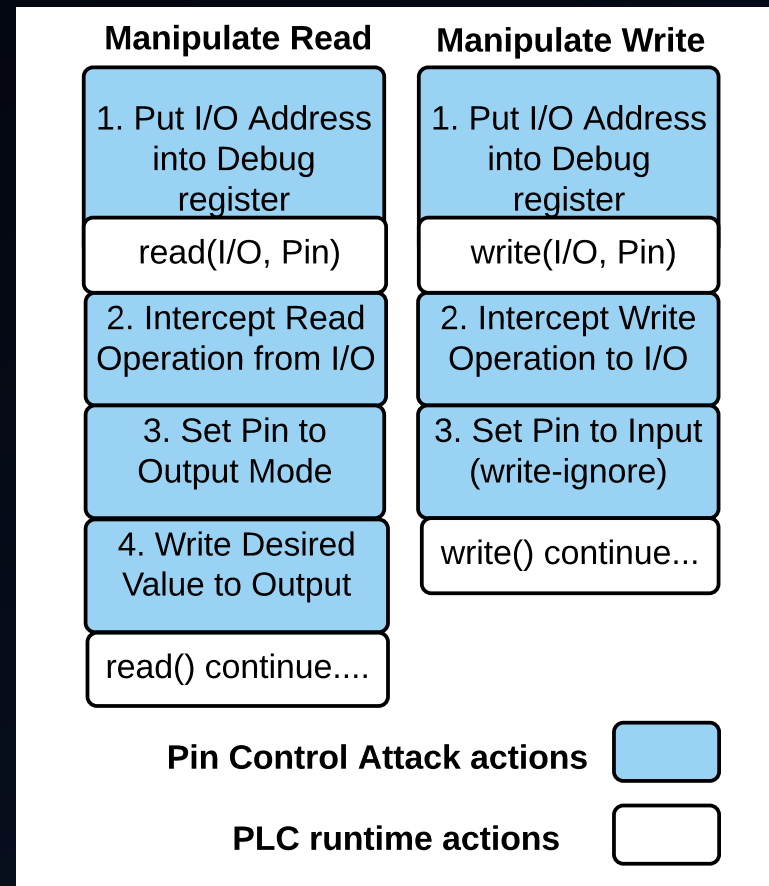
- Root privilege
- Knowledge of SoC registers
- Knowledge of mapping between I/O pins and the logic

2 Second version: C-code (shell code)

- Equal privilege as PLC runtime
- Knowledge of mapping between I/O pins and the logic

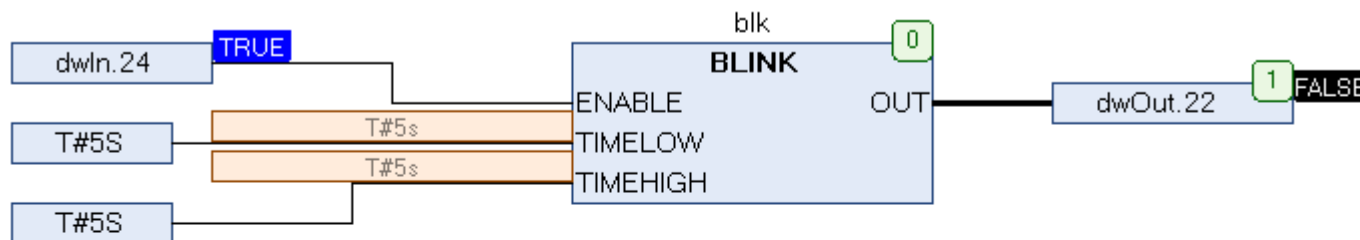
- No function hooking
- No modification of PLC runtime executable content
- No change to logic file

How Pin Configuration Attack Works?



Simple Logic

Lets test it with a simple Function Block
Language Logic.



Parameter	Type	Current ...	Pr...	Value	Default ...	Unit	Description
GPIO4	Enumeration of BYTE	not used		not used	not used		configuration of GPIO4
GPIO17	Enumeration of BYTE	not used		not used	not used		configuration of GPIO17
GPIO18	Enumeration of BYTE	not used		not used	not used		configuration of GPIO18
GPIO22	Enumeration of BYTE	Output		Output	not used		configuration of GPIO22
GPIO23	Enumeration of BYTE	not used		not used	not used		configuration of GPIO23
GPIO24	Enumeration of BYTE	Input		Input	not used		configuration of GPIO24
GPIO25	Enumeration of BYTE	not used		not used	not used		configuration of GPIO25

input : State of *In.24*
output: State of *Out.22*

Main Logic;

while *True* **do**

 read input;

while *input True* **do**

 switch_state(output, five seconds);
 //states are High or Low.

end

if *input False* **then**

 hold the state of the output;

else

 go to first while;

end

end

Simple Logic 2

- Second Logic for a real PLC

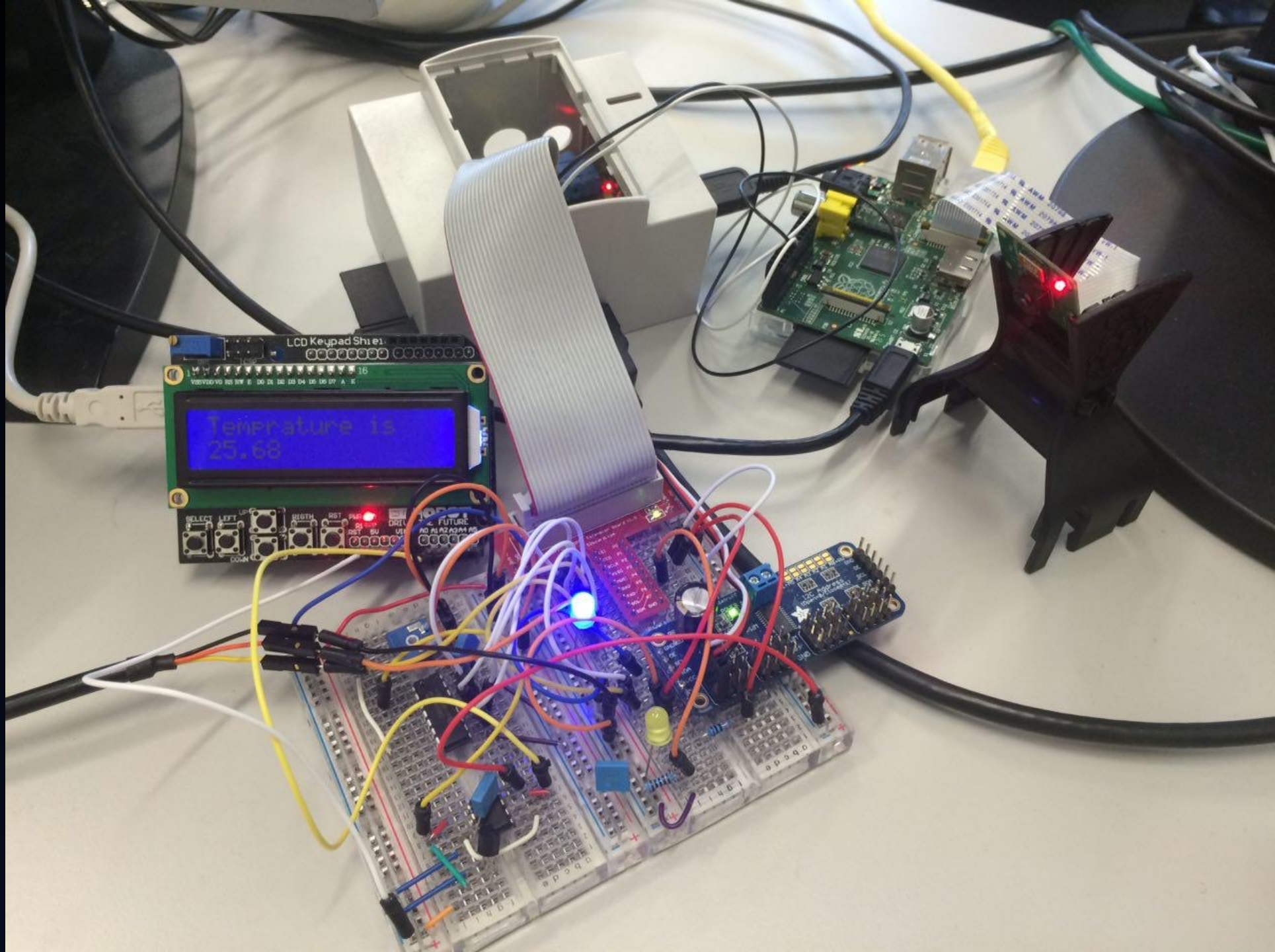
The screenshot displays the Siemens SIMATIC Manager interface. On the left, the 'Program structure' tree shows the project hierarchy: Applications > Application (1) > Library Manager > PLC_PRG > Task Configuration > PLC_Task (1). The main editor shows two programs for 'PLC_PRG':

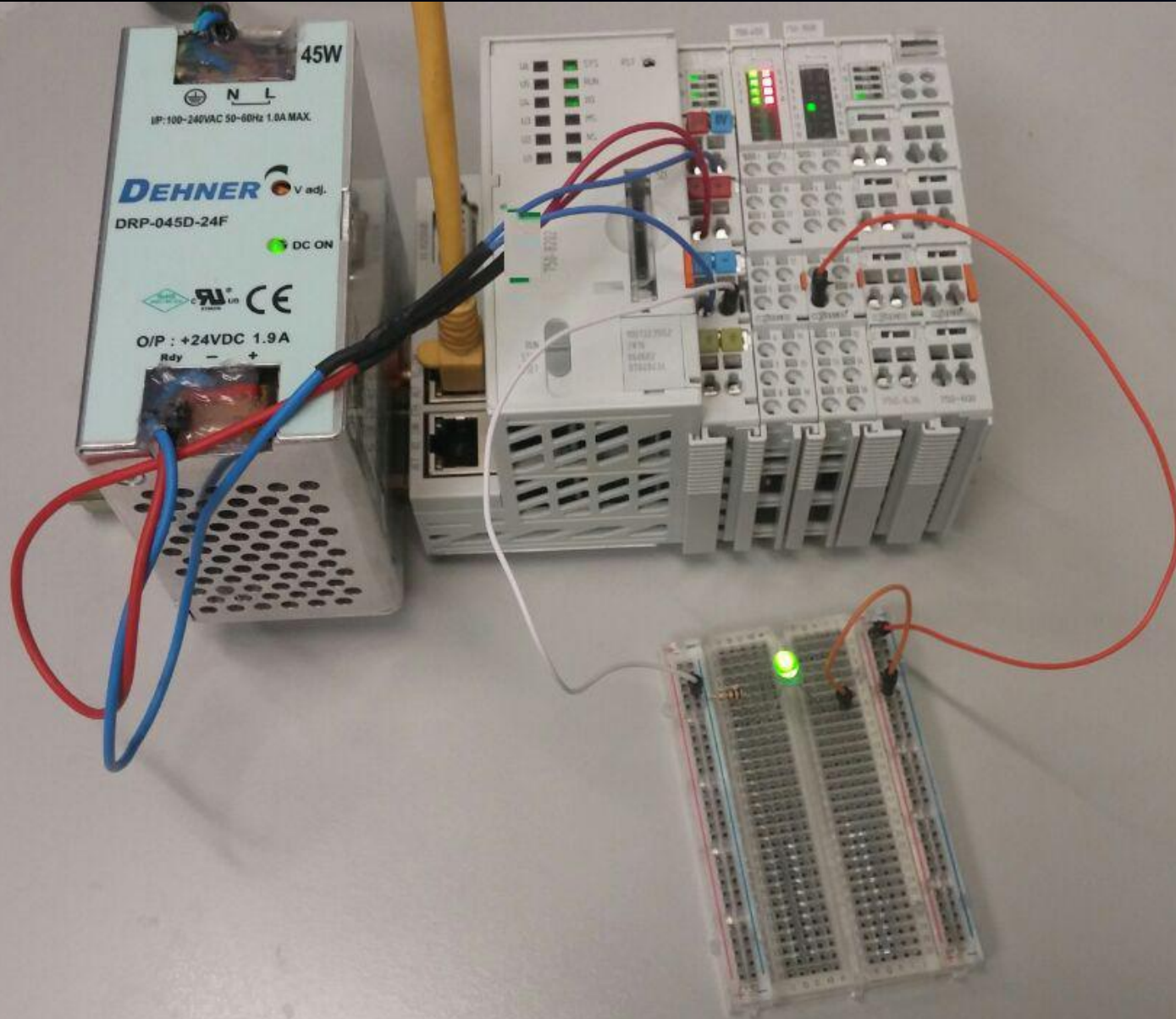
```
1 PROGRAM PLC_PRG
2 VAR
3   counter: INT := 0;
4   led_on_off: BOOL := FALSE;
5 END_VAR

1 // Simple on/off led switch
2 // Every 100 scan cycles
3 counter := counter + 1;
4
5 IF (counter = 100) THEN
6   counter := 0;
7   led_on_off := NOT led_on_off;
8 END_IF
```

On the right, a rack diagram shows the PLC hardware with modules U1 through U7. Below it, the 'I/O Mapping - 8_DO_DI_Generic' table is visible:

Mapping	Channel	Address	Type	Default Value	Current Value	Prepared Value	Unit	Description
	_IN	%IB 10	BYTE		0			Input Channels
	_OUT	%QB0	BYTE					Output Channels
Application.PLC_PRG.led_on_off	_OUT	%QX0.0	BOOL	FALSE	TRUE			Digital output
	_OUT	%QX0.1	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.2	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.3	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.4	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.5	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.6	BOOL	FALSE	FALSE			Digital output
	_OUT	%QX0.7	BOOL	FALSE	FALSE			Digital output





Lets look at it.

Demo 2

Digital

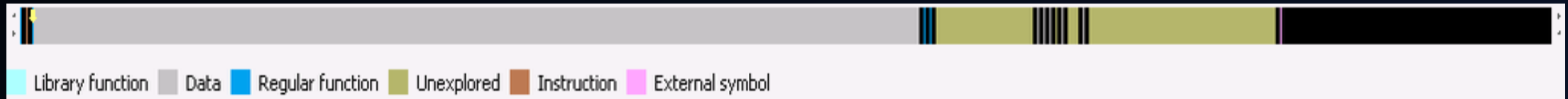
Lets look at it.

Demo 3

Digital

A PLC runtime Dynamic and Static Analysis

- I/O Mapping
- Look for Base Addresses of I/O

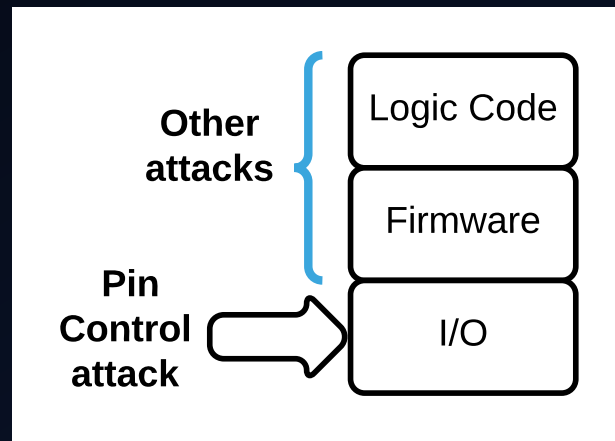


```
[b6e47f54] open("/etc/3S.dat", O_RDONLY) = 8 <0.001979>
[b6df334c] close(8) = 0 <0.001878>
[b6e47f54] open("/proc/cpuinfo", O_RDONLY) = 8 <0.001354>
[b6df334c] close(8) = 0 <0.007677>
[b6f2c7e4] open("/dev/mem", O_RDWR) = 8 <0.001182>
[b6e53998] mmap2(NULL, 4096, PROT_READ|PROT_WRITE, MAP_SHARED, 8, 0x2020)
[b6f2b000] close(8) = 0 <0.001246>
[b6f2c7e4] open("/dev/mem", O_RDWR) = 8 <0.001340>
[b6f2c7e4] open("/dev/spidev0.0", O_RDWR) = 9 <0.001886>
[b6e4fadc] ioctl(9, 0x80016b01, 0xb6a8d714) = 0 <0.001888>
```

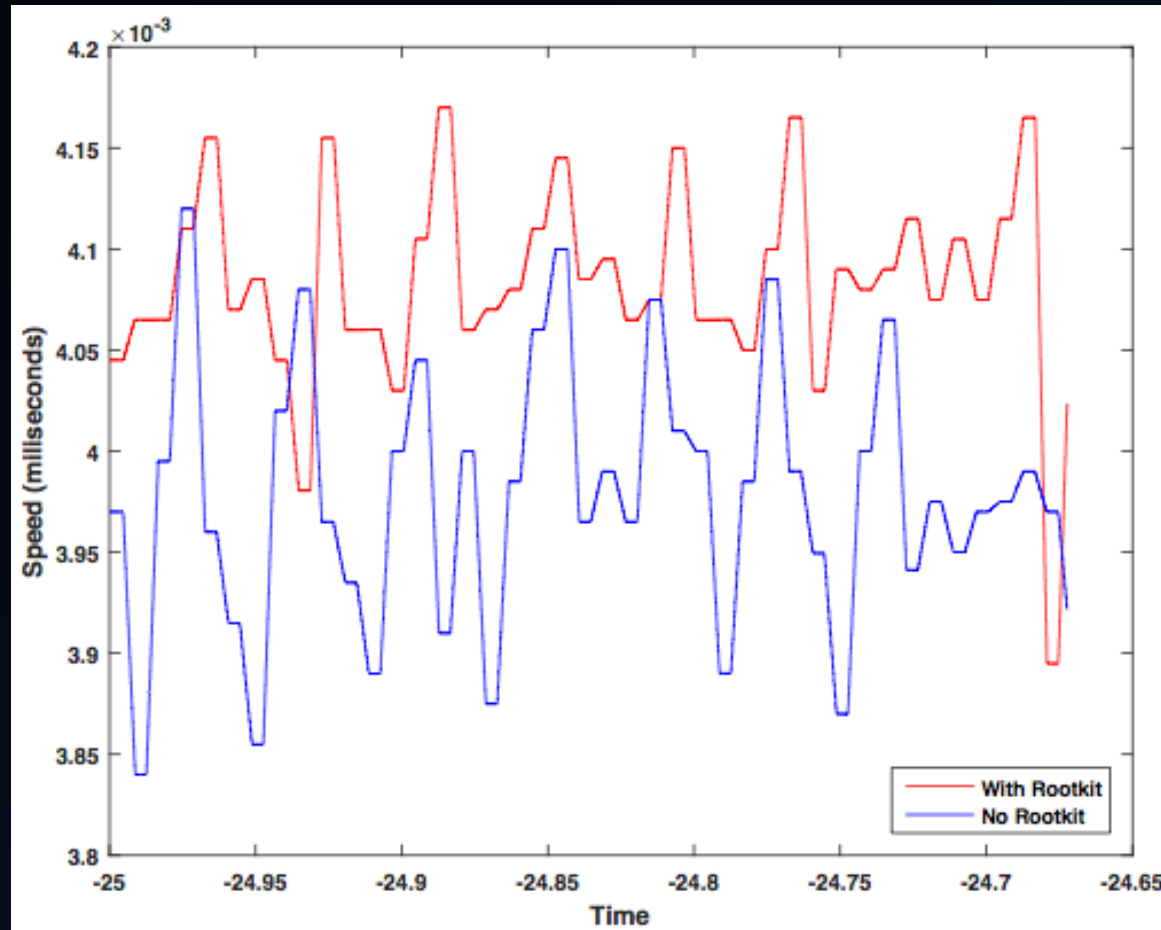
DCD	0x72322A06	0x8003FB5D	0xD182CBBC	0x575B9332	0x836A03F9
DCD	0xD1F2679A	0xD1B89713	0xD9BAA704	0xB6C6F738	0x3FA85B8C
DCD	0xB78538F5	0x3D4C2D10	0x282FF5B0	0x2B081994	0x1848E56D
DCD	0xAA8C7B82	0xDE23AF80	0xE6144FFD	0xFE82BA1B	0x18604BA9
DCD	0x223D3D45	0x1A00B106	0x825AC9E5	0xF425FFE6	0xB19B375B
DCD	0xEF878EA7	0x172C1C83	0x40E54D04	0x588CDBC8	0x1B19AC0F
DCD	0x7ED50852	0xE0C950C8	0x9C67C354	0x3DA8F807	0x421FBB11
DCD	0x2C816A17	0xFB3E4375	0x60DB663	0xF15C6122	0x64514032
DCD	0xDA75AF94	0x9929D1B3	0x3D910885	0x8984059F	0x4F66A58
DCD	0x97F2C7D9	0x4808685D	0xB24602AE	0x75828FCA	0x734C7E16
DCD	0xD39E5C65	0x4BF3B903	0x952C30A7	0x2F92553B	0xD2FBF6E7
DCD	0xD2AD2C07	0x47B449B9	0x46CF816A	0x5B1A9B0D	0x9B61780
DCD	0xE7864462	0xA5DA033E	0x4B3A8C38	0xA57A4FD0	0x235575B9
DCD	0xB2E7AEC8	0xCC77010F	0xD6A729C	0x8BF267AC	0xB91822D4
DCD	0x50225EB2	0x670E0020	0x05EE2700	0x0210200E	0x00E01E70

I/O Attack: Rootkit

- Rootkit needs root user to install its code as a Loadable Kernel Module (LKM).
- vmalloc() allocates our LKM. It evades Doppelganger.
- Do not do any kind of function hooking, evades Autoscopy Jr.
- Can change the logic regardless of logic operation.



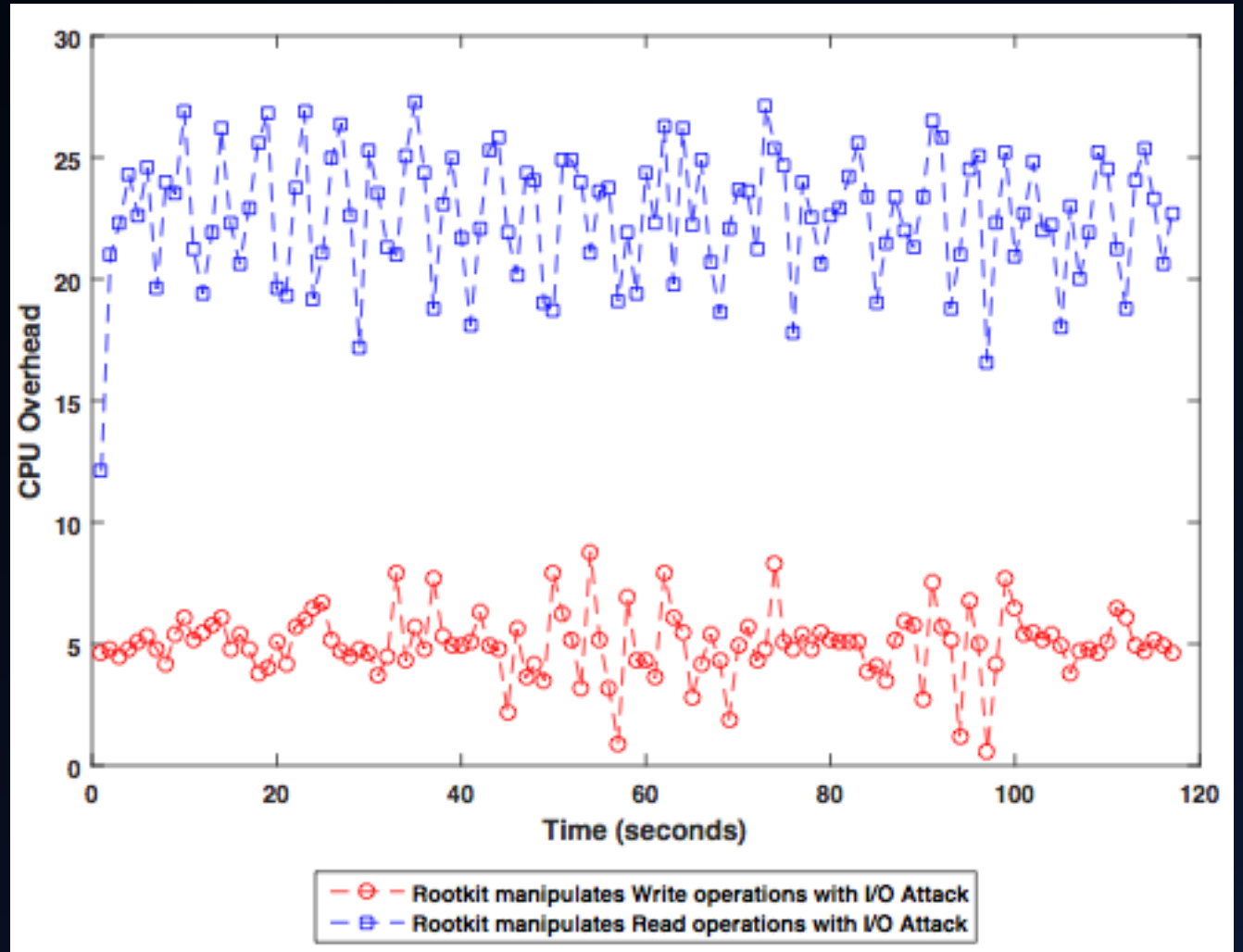
I/O response time fluctuation in rootkit variant



CPU Overhead

Write Manipulation: ~ 5%

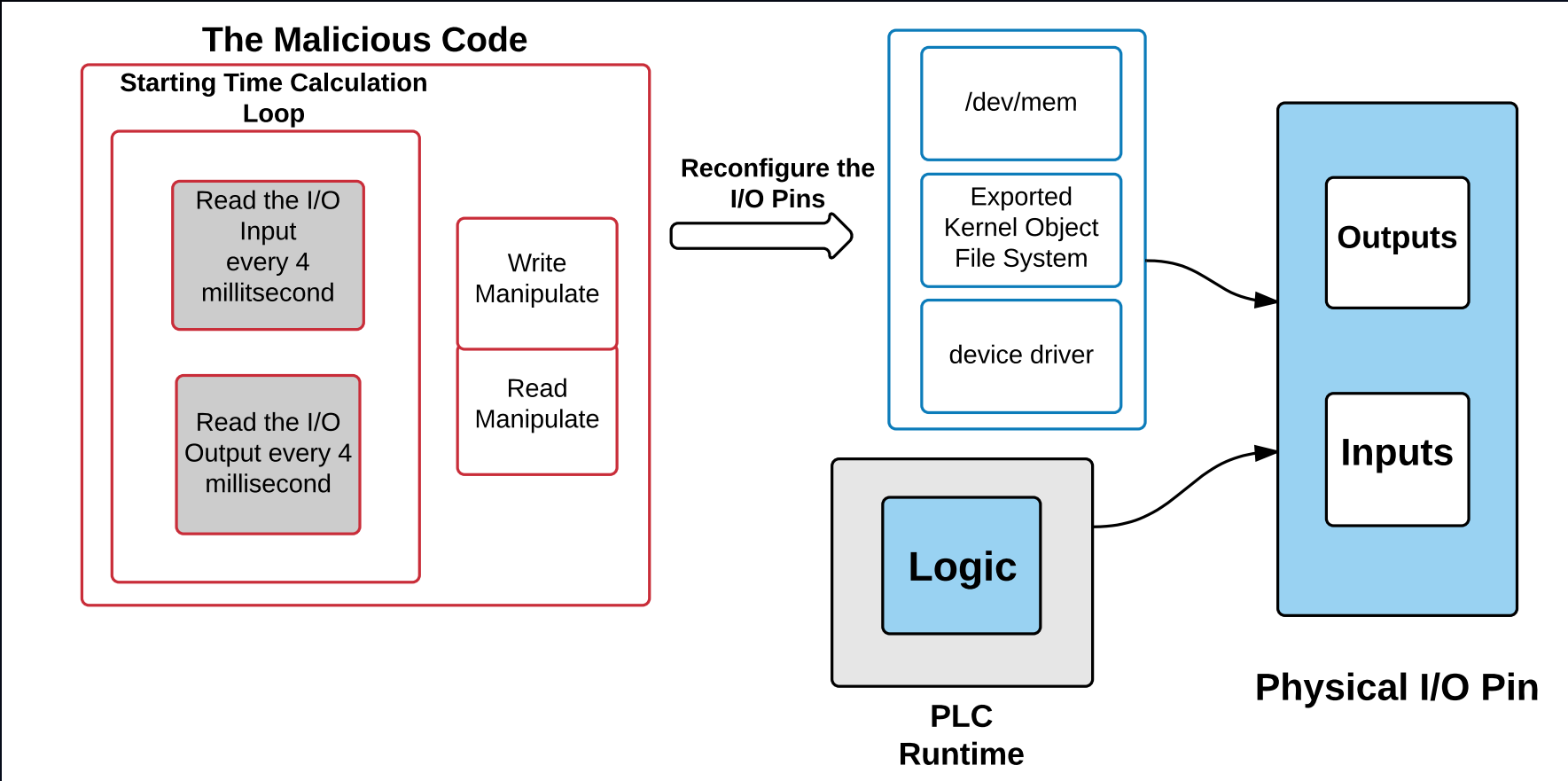
Read Manipulation: ~ 23%



Second Variant of the Attack – No Rootkit !

- No need to have rootkit!
- We can do the same with the PLC runtime privilege.
- Overhead below 1%.
- We can either remap the I/O or use already mapped I/O address.
- As shellcode

Second variant



Second Variant

Manipulate Read

1. Find the Reference Starting Time

3. Set Pin to Output Mode (write-enable)

4. Write Desired Value to Output Pin

read(I/O, Pin)

Manipulate Write

1. Find the Reference Starting Time

3. Set Pin to Input (write-ignore)

write() to I/O

3. Set Pin to Output (write-enable)

Write desired value

Pin Control Attack actions



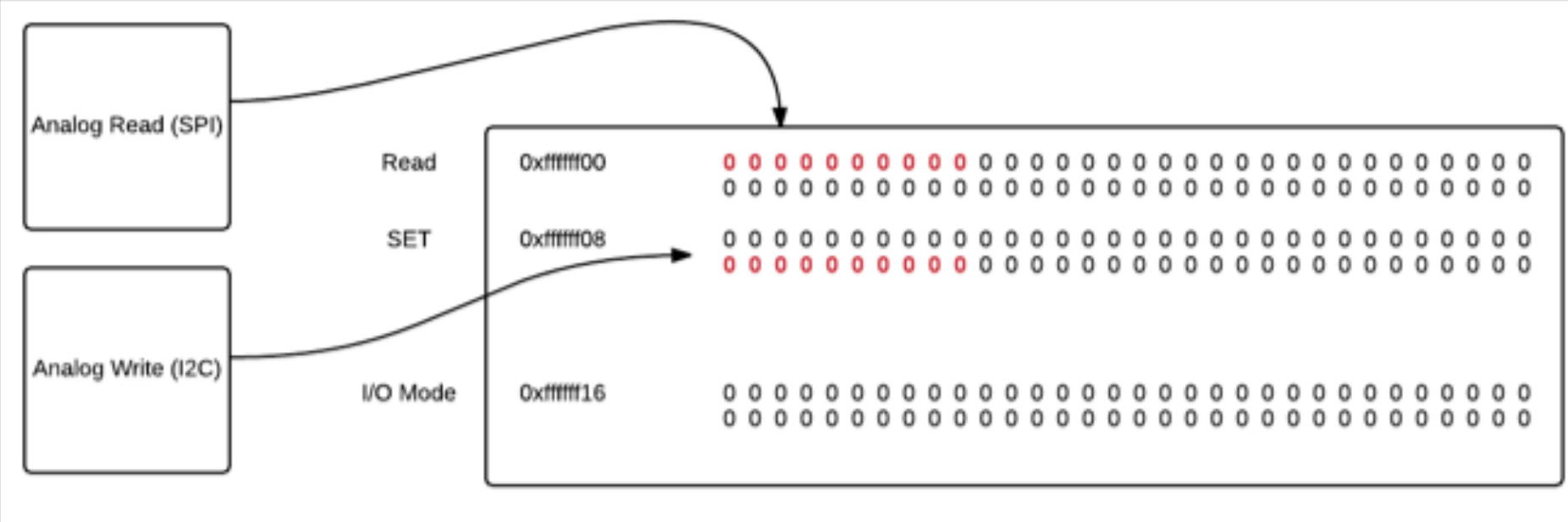
PLC runtime actions



What about Analog Control?

- Analog signals are basically aggregation of digital signals.
- Two ways to do it:
 - 1. If part of or entire analog memory can get multiplexed to digital pins attacker can multiplex the pin and write digital bits and basically control the values in the analog memory
 - 2. Using the technique which we can PC+1, we tell the interrupt handler to return the control to the next instruction within the PLC runtime, basically avoiding write operation occur

Analog I/O Manipulation



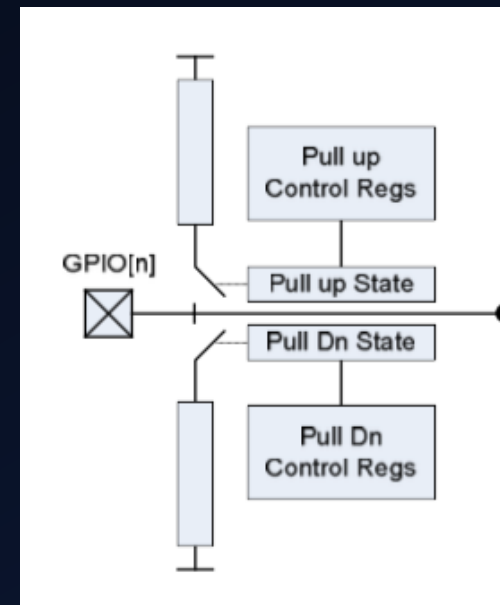
Lets look at it.

Demo

Analog

Other Future Possibilities!

- Attacking pull-up and pull-down resistors in I/O interfaces
- What if we disable them?
- Remotely manipulate the I/O via a powerful electromagnetic field!





Never trust your inputs!

Discussions

- For now attacker can:
 - Simply change the logic
 - Modify PLC Runtime executable
- Fixing these attacks are trivial:
 - Proper Authentication
 - Proper Logic Checksum
 - PLC Runtime integrity verification
- Next Step for attackers:
 - Achieve its goal without actually modifying the Logic or Runtime or hooking functions

Race to the Bottom

**RACE TO THE
BOTTOM**



As soon as security is introduced at some layer of computer or network architecture abstraction, the attackers are going one layer down.

In the hacking community it is called Race-to-the-Bottom

Conclusions

- Need to focus on system level security of control devices In future more sophisticated techniques come that evade defenses.
 - Pin Control attack is an example of such attacks.
- Pin Control Attack:
 - lack of interrupt for I/O configuration registers
 - Significant consequences on protected PLCs and other control devices such as IEDs.
- Solution:
 - It is hard to handle I/O interrupts with existing real-time constraints.
 - Monitoring I/O Configuration Pins for anomalies.
 - User/Kernel space separation for I/O memory.

Questions?


Looking for more...

Attend our talk at DigitalBond S4x17, Miami, USA

Everything that has a beginning has an end.

The Matrix Revolutions.

Contact:

a.abbasi@utwente.nl  @bl4ckic3

 @m4ji_d