

CCC '07

Karsten Nohl, Starbug, Henryk Plötz

# MIFARE SECURITY

# RFID tags

- ◆ Radio Frequency IDentification
- ◆ Tiny computer chips
- ◆ Passively Powered



# RFID Applications

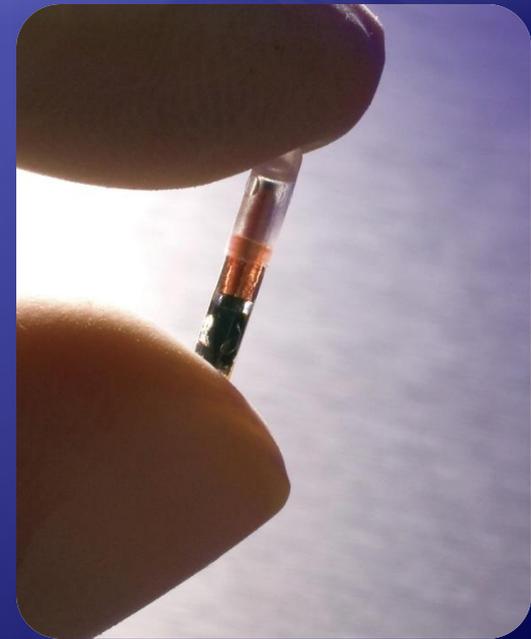
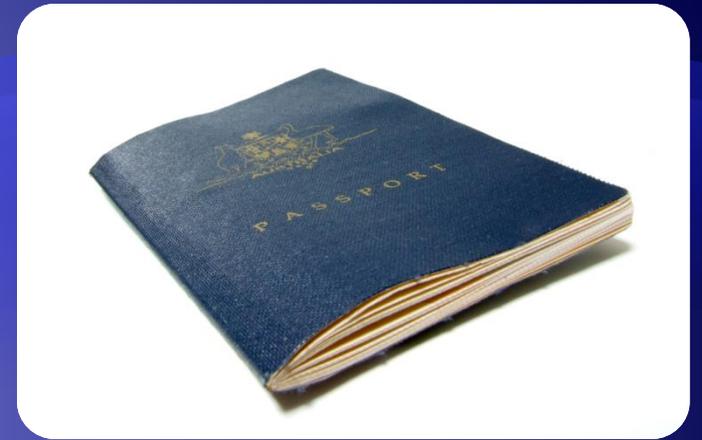
- ◆ RFIDs become ubiquitous
- ◆ Integrated in many *security* applications
  - ◆ Tickets
  - ◆ Access Control
  - ◆ Car Ignition



# RFID Trends

- ◆ Passports
- ◆ Implants
- ◆ ...

RFIDs become *universal identifier*. Might replace passwords, PINs, and fingerprints.



# RFID Trends (II)

- ◆ Tagging of consumer goods
  - ◆ Will replace bar-codes!
- ◆ Threat to *Privacy*
  - ◆ Customer tracking
  - ◆ Leaks internal business information!



# RFID Tre

- ◆ Tagging of c
  - ◆ Will replace
- ◆ Threat to *Pr*
  - ◆ Customer t
  - ◆ Leaks inter



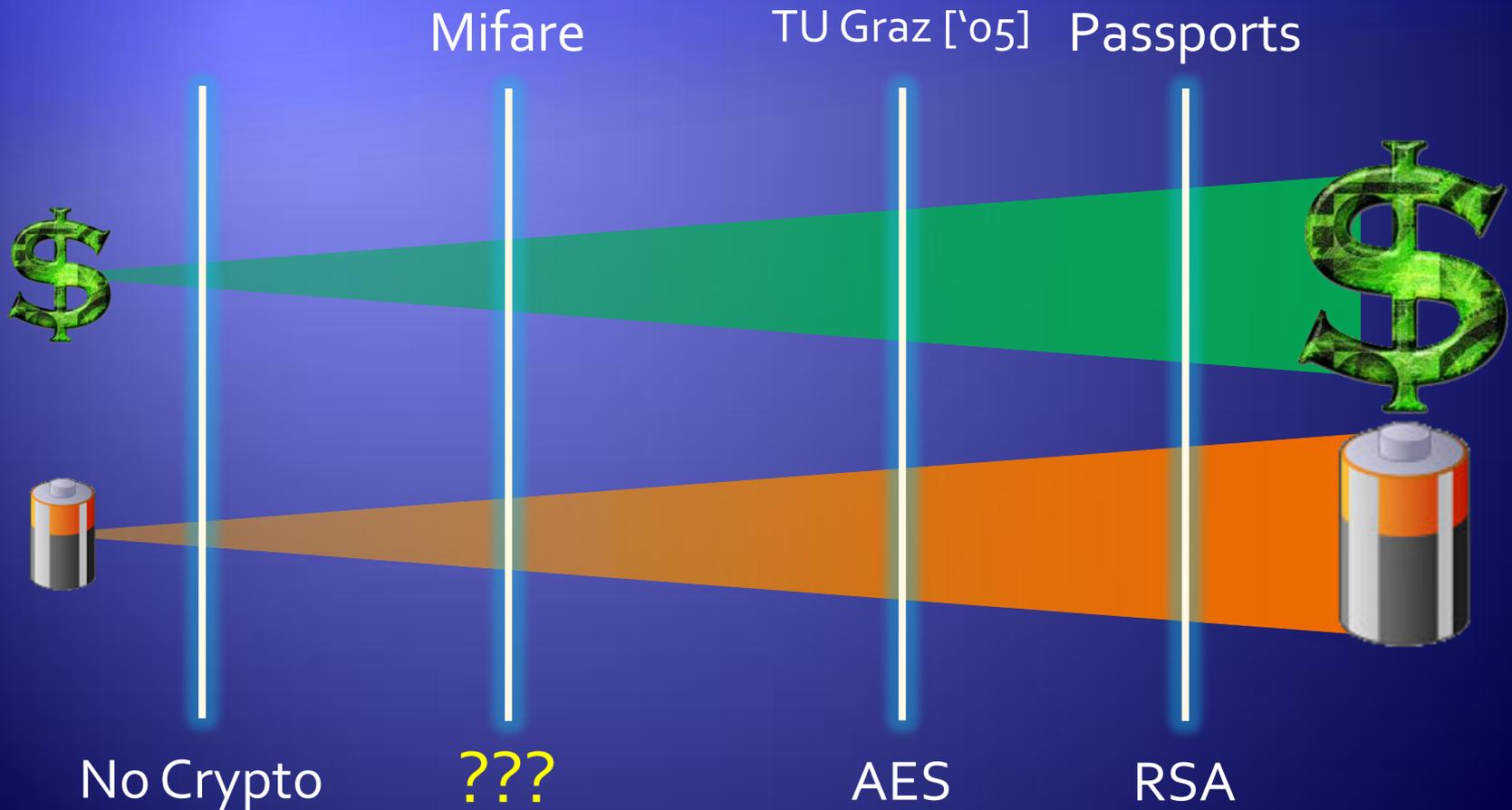
# Motivation

*Cryptography* on RFIDs is needed for:

- ◆ Unclonability
  - ◆ Credit cards, luxury goods, medication, ...
- ◆ Privacy!

But, what crypto is small enough for tags?

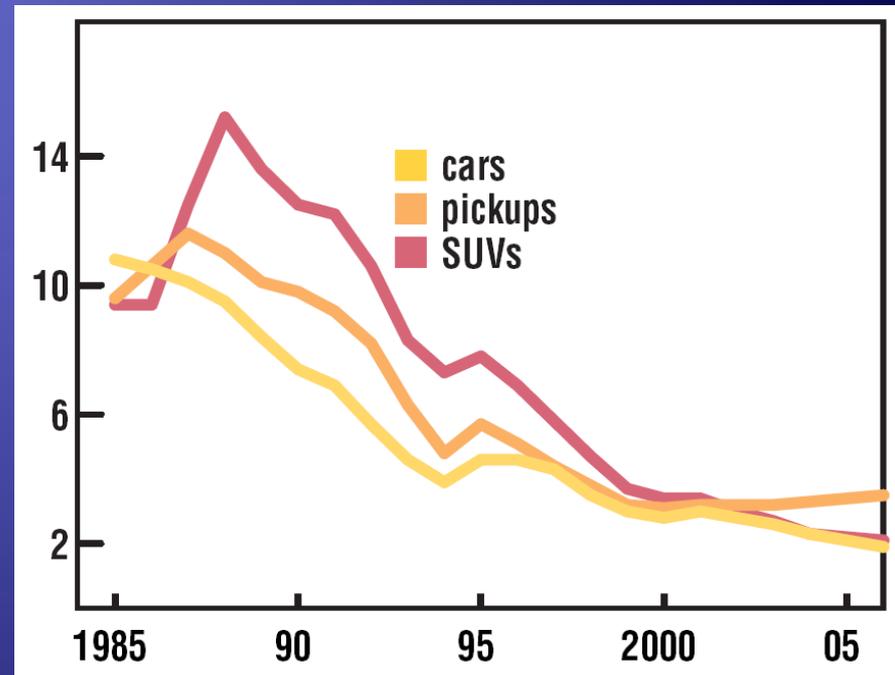
# RFID-Crypto Mismatch



# Mifare Security

- ◆ Philips claims:
  - ◆ “approved authentication”
  - ◆ “advanced security levels”
- ◆ 48 bit key

Car thefts  
(source: hldi.org)



# Our Project

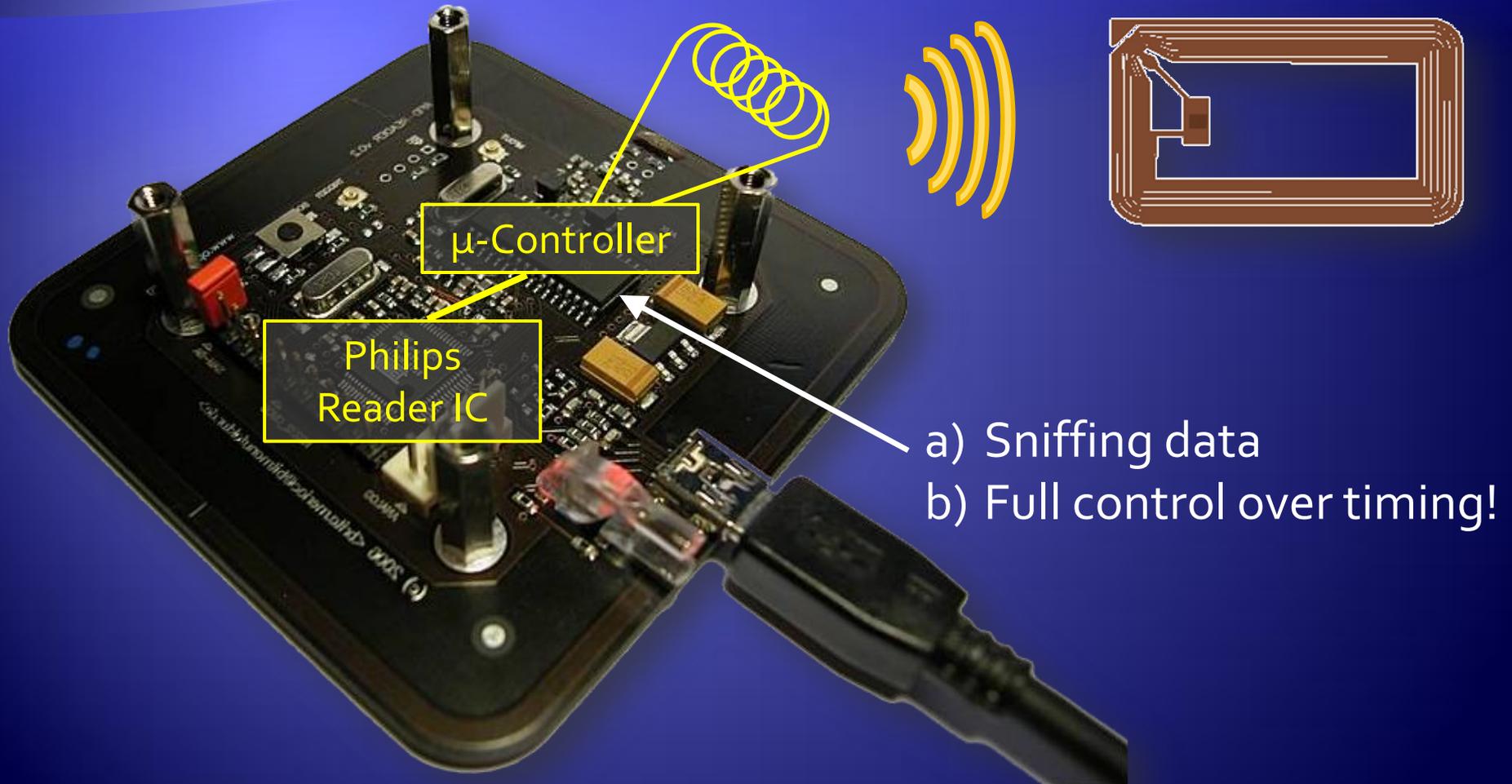
Reverse-engineering of the Mifare crypto and evaluating its security

- ◆ Reconstruct circuit from photos of chip
- ◆ Sniff reader-tag communication

verify



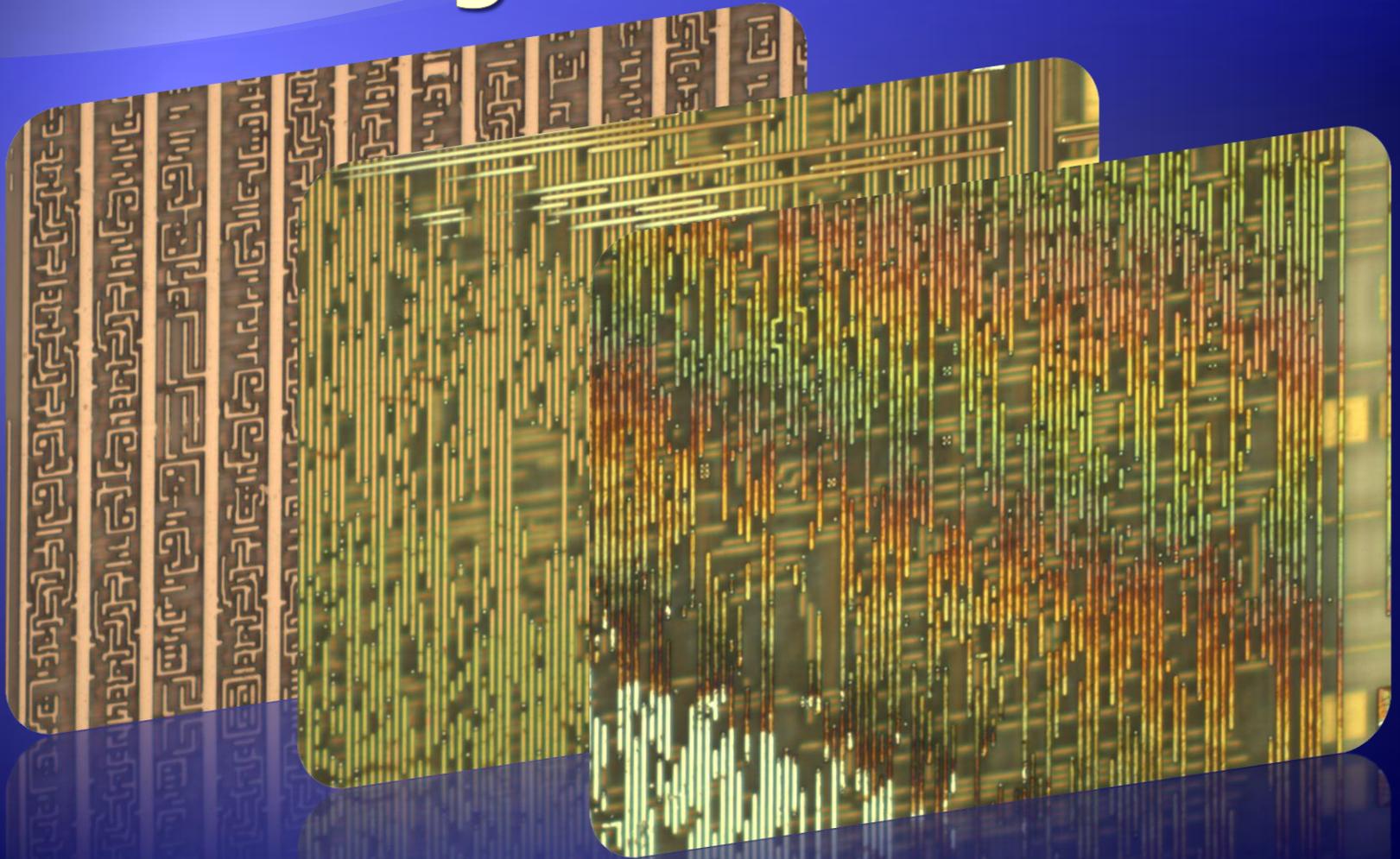
# Hardware: OpenPCD (+PICCC)



# Mifare RFID tag

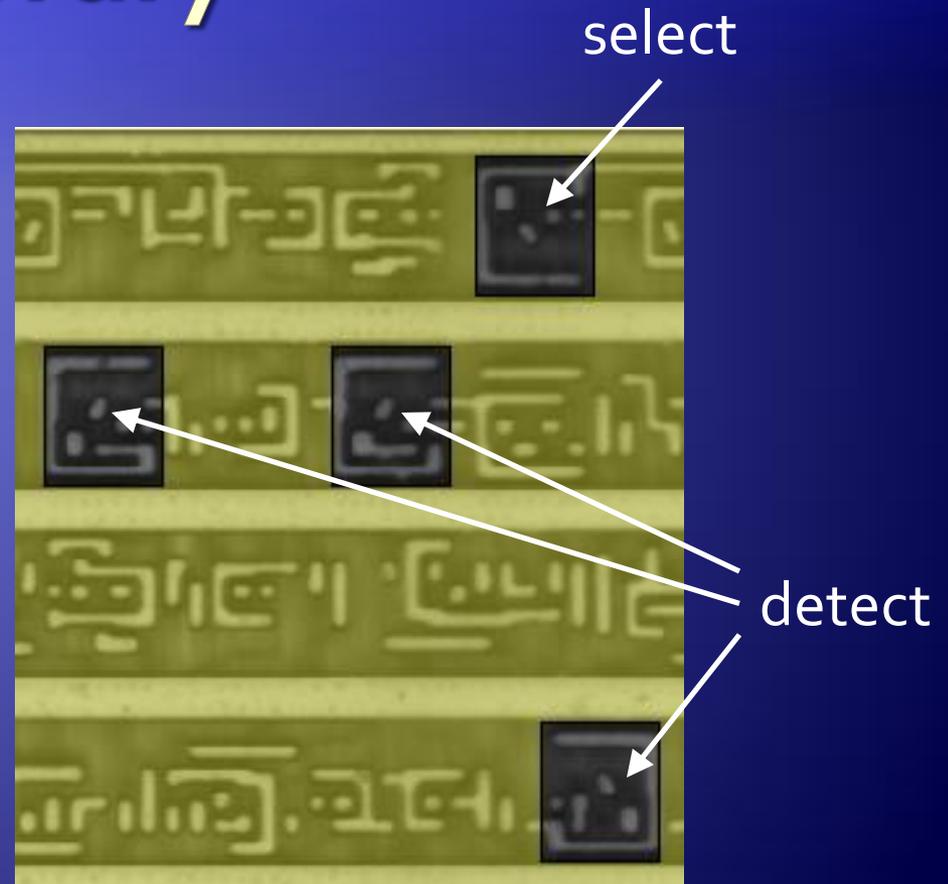


# Circuit Images



# Step 1: build library

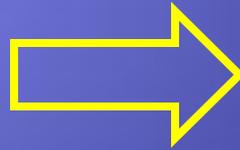
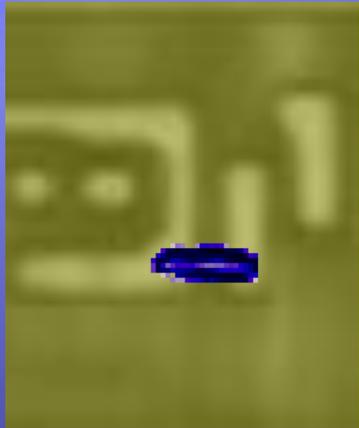
- ◆ Chip has several thousand gates
- ◆ But only ~70 different types
  - ◆ Detection can be automated



# Step 2: Encircle Crypto

- ◆ Even tiny RFID chip too large to analyze entirely
  - ◆ Crypto <10% of gates!
- ◆ Focus on interesting-looking parts:
  - ◆ Strings of flip-flops (registers)
  - ◆ XOR
  - ◆ Units around edges that sparsely connected to the rest of the chip

# Step 3: Reproduce Circuit

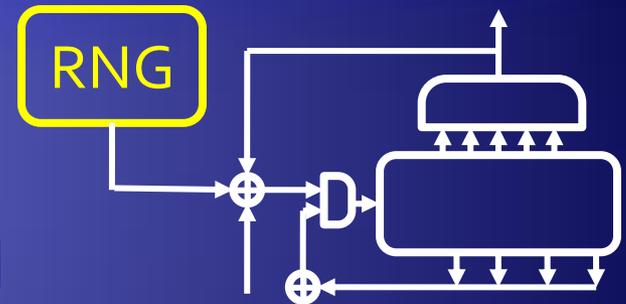


- ◆ Very error-prone and tedious process
  - ◆ Will automate further



# Random Number Generator

- ◆ 16(!!)-bit random numbers
  - ◆ LFSR –based
  - ◆ Value derived from time of read



Our Attack:

- ◆ Control timing (OpenPCD)
  - = control random number (works for tag and reader!)
  - = break Mifare security :)



WIKIPEDIA  
The Free Encyclopedia

#### navigation

- [Main Page](#)
- [Contents](#)
- [Featured content](#)
- [Current events](#)
- [Random article](#)

#### interaction

- [About Wikipedia](#)
- [Community portal](#)
- [Recent changes](#)
- [Contact Wikipedia](#)
- [Donate to Wikipedia](#)
- [Help](#)

#### search

#### toolbox

- [What links here](#)
- [Related changes](#)
- [Upload file](#)
- [Special pages](#)

[article](#)[discussion](#)[edit this page](#)[history](#)

# Linear feedback shift register

From Wikipedia, the free encyclopedia

(Redirected from [LFSR](#))

A **linear feedback shift register** (LFSR) is a [shift register](#) whose input bit is a [linear function](#) of its previous

The only linear functions of single bits are xor and inverse-xor; thus it is a shift register whose input bit is driven

The initial value of the LFSR is called the seed, and because the operation of the register is deterministic, the state. Likewise, because the register has a finite number of possible states, it must eventually enter a repeating cycle which appears random and which has a very long cycle.

Applications of LFSRs include generating [pseudo-random numbers](#), [pseudo-noise sequences](#), fast digital counting, and are very common.

## Fibonacci LFSRs

The list of the bits' positions that affect the next state is called the tap sequence. In the diagram below, the sequence of bits is shifted to the right, the output bit is taken from the rightmost bit, and then fed back into the leftmost bit.

- The outputs that influence the input are called *taps* (blue in the diagram below).
- A maximal LFSR produces an *n-sequence* (i.e. cycles through all possible  $2^n - 1$  states within the shift register) and never change.

The sequence of numbers generated by an LFSR can be considered a [binary numeral system](#) just as valid as any other

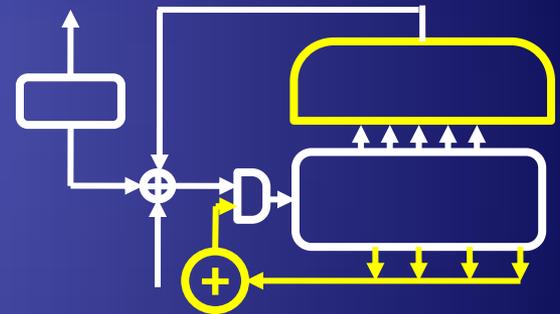
The tap sequence of an LFSR can be represented as a [polynomial mod 2](#). This means that the coefficients of the polynomial are either 0 or 1. For example, if the taps are at the 16th, 14th, 13th and 11th bits (as below), the resulting LFSR polynomial is

$$x^{16} + x^{14} + x^{13} + x^{11} + 1$$

The 'one' in the polynomial does not correspond to a tap - it corresponds to the input to the first bit (i.e.  $x^0$ , where  $x^0 = 1$ ). The first and last bits are always connected as an input and tap respectively.

# Structural Weaknesses

- 1) No non-linear component in feedback loop  
→ No forward secrecy
- 2) Output bit derived from fixed subset of bits  
→ non-optimal avalanche properties



Suggests attack on key faster than brute-force (known-plaintext)

# Brute-Force Attack

- ◆ Cipher complexity low
  - ◆ Has probably been the highest design goal
  - ◆ Allows for very efficient FPGA implementation

\$100 key cracker will find key  
in ~1 week! (much faster even  
when trading space for time)



# Mifare Security

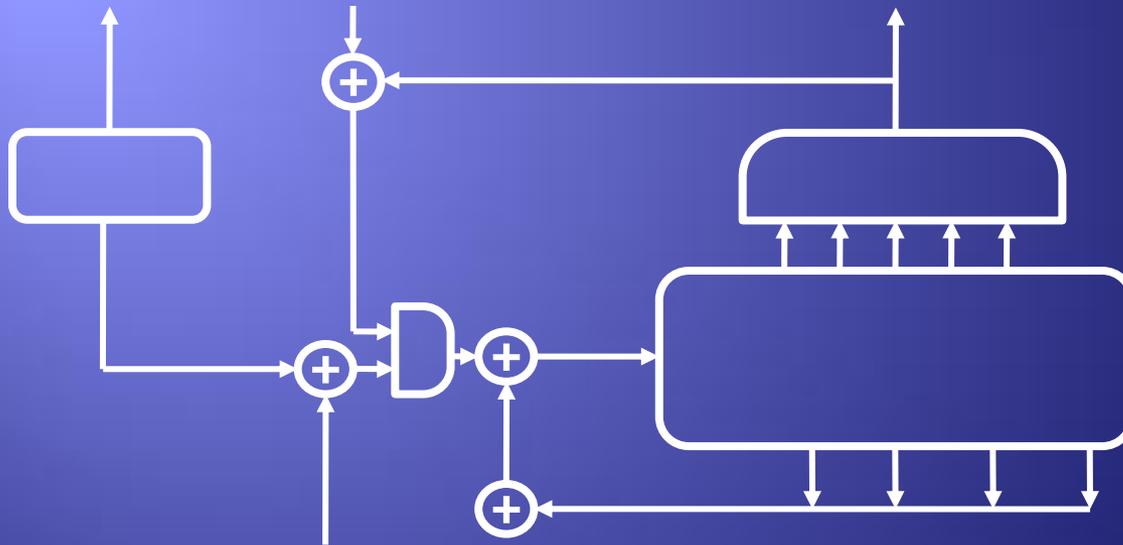


- ◆ Protection perhaps sufficient to protect transactions of very small value
  - ◆ E.g., Micro-payments, privacy
- ◆ Security too weak for:
  - ◆ Access control, car theft protection, credit cards, ...

# Lessons Learned

- ◆ Obscurity and proprietary crypto add security only in the short-run
  - ◆ (but lack of peer-review hurts later)
- ◆ Constraints of RFIDs make good crypto extremely hard
  - ◆ Where are the best trade-offs?
  - ◆ How much security is needed?

# Questions?



Karsten Nohl  
nohl@virginia.edu